

Theory of Robotics Systems Graph Search





Faculty of Technical Sciences



Ss. Cyril and Methodius University Faculty of Electrical Engineering and Information Technologies

Zagreb University of Applied Sciences



School of Electrical Engineering University of Belgrade



Faculty of Physics Warsaw University of Technology



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Metode pretrage grafova







Primer 1

- Putanja od zelene ćelije do crvene
- 8 koraka iz jedne ćelije
- Koja je najkraća putanja?
- Šta ako imamo 4 moguća

koraka iz ćelije.



Assumptions about robot motion





Primer 1

- Pretražimo dobijeni graf.
- Pretpostavljamo da sve ivice grafa mogu kreirati putanju.
- Rešimo problem tako što nađemo putanju od A do B.







Pretraga grafova

- Ivice grafa mogu imati pridružene **težine**
 - Cena putanje
- Ima smisla da se razmatraju samo pozitivne težine.
- Algoritmi sa minimalnom cenom nikad ne posećuju isti čvor dva puta.
- Najpoznatije pretrage:
 - Pretraga u širinu (BFS)
 - Pretraga u dubinu (DFS)
 - Dijkstra's Algoritam
 - A*



Directed graph: edges Weighted graph: edges

edges have direction edges have costs





Drvo pretrage

• Konstruišemo "drvo" pomoću kojeg vršimo pretragu.





Drvo pretrage

- Drvo pretrage:
 - Početno stanje se uzima kao koren
 - Naslednici skup svih mogućih akcija
 - Uglavnom želimo da izbegnemo konstrukciju kompletnog stabla







Osnove pravolinijske pretrage

- Počnemo od korena (početno stanje), širimo drvo sve dok ne dođemo do cilja.
- Proširivanje nodova znači dodavanje njihovih naslednika.
- Treba probati da se nađe rešenje sa što manje proširivanja.
- Otvoren set čvorovi iz kojih se mogu vršiti proširivanja
- Zatvoren set svi oni čvorovi koje je algoritam posetio





Osnove pravolinijske pretrage

FOR	WARD_SEARCH
1	$Q.Insert(x_I)$ and mark x_I as visited
2	while Q not empty do
3	$x \leftarrow Q.GetFirst()$
4	if $x \in X_G$
5	return SUCCESS
6	forall $u \in U(x)$
7	$x' \leftarrow f(x, u)$
8	if x' not visited
9	Mark x' as visited
10	Q.Insert(x')
11	else
12	Resolve duplicate x'
13	return FAILURE





Pretraga po širini

FOR	WARD SEARCH
1	$Q.Insert(x_I)$ and mark x_I as visited
2	while Q not empty do
3	$x \leftarrow Q.GetFirst()$
4	if $x \in X_G$
5	return SUCCESS
6	forall $u \in U(x)$
7	$x' \leftarrow f(x, u)$
8	if x' not visited
9	Mark x' as visited
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11	else
12	Resolve duplicate x'
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Pretraga po širini

WARD_SEARCH
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Mark x' as visited
Q.Insert(x')
else
Resolve duplicate x'
return FAILURE

Open (Q):	Closed:
{A,C,G}	{B,A,C,G}







Pretraga po širini

FOR	WARD_SEARCH
2	while Q not empty do
3	$x \leftarrow Q.GetFirst()$
4	if $x \in X_G$
5	return SUCCESS
6	forall $u \in U(x)$
7	$x' \leftarrow f(x, u)$
8	if x' not visited
9	Mark x' as visited
10	Q.Insert(x')
11	else
12	Resolve duplicate x'
13	return FAILURE

	~	B	
		C	G
\bigcirc			



Open (Q):	Closed:
{C,G,D}	{B,A,C,G,D}











Open (Q):	Closed:
{G,D,I}	{B,A,C,G,D,I}





Pretraga po širini











Pretraga po širini





Open (Q): Closed: {I,J} {B,A,C,G,D,I,J}





Pretraga po širini











Pretraga po širini







Pretraga po dubini

FOR	WARD SEARCH
1	$Q.Insert(x_I)$ and mark x_I as visited
2	while Q not empty do
3	$x \leftarrow Q.GetFirst()$
4	if $x \in X_G$
5	return SUCCESS
6	forall $u \in U(x)$
7	$x' \leftarrow f(x, u)$
8	if x' not visited
9	Mark x' as visited
10	Q.Insert(x')
11	else
12	Resolve duplicate x'
13	return FAILURE

C

Our (DFS) queue will be LIFO:

Closed:

{}

Open (Q):

{B}

- push (Q.Insert) onto the front
- pop (*Q.GetFirst*) from the front





Pretraga po dubini



Open (Q): Closed: {A,C,G} {B,A,C,G}

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G



Pretraga po dubini

FOR 1	WARD_SEARCH $Q.Insert(x_I)$ and mark x_I as visited
2	while Q not empty do
3	$x \leftarrow Q.GetFirst()$
4	if $x \in X_G$
5	return SUCCESS
6	forall $u \in U(x)$
7	$x' \leftarrow f(x, u)$
8	if x' not visited
9	Mark x' as visited
10	Q.Insert(x')
11	else
12	Resolve duplicate x'
13	return FAILURE





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Open (Q): Closed: {D,C,G} {B,A,C,G,D}



Pretraga po dubini





Open (Q): {I,J,C,G} Closed: {B,A,C,G,D,I,J}





Pretraga po dubini





Open (Q): {J,C,G} Closed: {B,A,C,G,D,I,J}





Pretraga po dubini

FOR	WARD_SEARCH		
1	$Q.Insert(x_I)$ and ma	$\operatorname{ark} x_I$ as visited	1
2	while Q not empty q	lo	
3	$x \leftarrow Q.GetFirst($)	
4	if $x \in X_G$		
5	return SUCC	ESS	
6	forall $u \in U(x)$		
7	$x' \leftarrow f(x, u)$		
8	if x' not visite	d	
9	Mark x' as	visited	
10	Q.Insert(x)	r')	
11	else		
12	Resolve du	plicate x'	
13	return FAILURE	Open (Q):	Closed:
		8	{B,A,C,G,D,I,J}









Uporedna analiza







DFS





Akcije sa cenom

- Kako rešiti slučaj ako određena putanja ima veću cenu nego neka druga?
- Najpopularniji algoritmi Dijkstra's i A*







Dijkstra's algoritam

- Objavljen od strane Edsger Dijsktra 1959.
- Otvara čvorove sa najmanjom cenom najbliže od korena grafa.
- Najkorišćeniji algoritam za rešavanje problema trgovačkog putnika.







Dijkstra's algoritam

WARD SEARCH		
$Q.Insert(x_l)$ and mark x_l as vis	ited	
while Q not empty do		
$x \leftarrow Q.GetFirst()$		
if $x \in X_G$		
return SUCCESS		
forall $u \in U(x)$		
$x' \leftarrow f(x, u)$		
if x' not visited	Open (Q):	Closed:
Mark x' as visited	{B(0)}	{B(0)}
Q.Insert(x')	1	
else		
Resolve duplicate x'	\	
return FAILURE		Diiketro que
	WARD SEARCH $Q.Insert(x_I)$ and mark x_I as vis while Q not empty do $x \leftarrow Q.GetFirst()$ if $x \in X_G$ return SUCCESS forall $u \in U(x)$ $x' \leftarrow f(x, u)$ if x' not visited Mark x' as visited Q.Insert(x') else Resolve duplicate x' return FAILURE	WARD SEARCH $Q.Insert(x_I)$ and mark x_I as visited while Q not empty do $x \leftarrow Q.GetFirst()$ if $x \in X_G$ return SUCCESS forall $u \in U(x)$ $x' \leftarrow f(x, u)$ if x' not visited Open (Q): Mark x' as visited {B(0)} Q.Insert(x') else Resolve duplicate x' return FAILURE



 $A = \begin{bmatrix} 5 & 2 & 2 & 5 \\ 2 & 2 & 5 \\ 3 & 2 & 1 \\ 3 & 2 & 1 \\ 0 & 2 \end{bmatrix}$

Our Dijkstra queue will be ordered by cost to arrive:

- push (*Q.Insert*) by cost
- pop (Q.GetFirst) from the front, and add it to the closed list





Dijkstra's algoritam

1	$Q.Insert(x_I)$ and mark x_I as visited
2	while Q not empty do
3	$x \leftarrow Q.GetFirst()$
4	if $x \in X_G$
5	return SUCCESS
6	forall $u \in U(x)$
7	$x' \leftarrow f(x, u)$
8	if x' not visited
9	Mark x' as visited
10	Q.Insert(x')
11	else
12	Resolve duplicate x'
13	return FAILURE



Closed:

{ B (0) }



Open (Q): { C (2), G (3), A (5) }





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Dijkstra's algoritam



A (4) }



Dijkstra's algoritam







Open (Q): { G (3), A (4), D (4) }

Closed: { B (0), C (2) }





Dijkstra's algoritam

FOR 1	WARD_SEARCH $Q.Insert(x_I)$ and mark x_I as visited	
2	while Q not empty do	
3	$x \leftarrow Q.GetFirst()$	
4	if $x \in X_G$	
5	return SUCCESS	
6	forall $u \in U(x)$	
7	$x' \leftarrow f(x, u)$	
8	if x' not visited	
9	Mark x' as visited	
10	Q.Insert(x')	
11	else	
12	Resolve duplicate x'	
13	return FAILURE	Open (





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5

3

2

n

2

G

2

























Dijkstra's algoritam





Open (Q): Closed: { I (5), { B (0), J (6) } C (2),

G (3), A (4), Co-funded by the D (4) } Erasmus+ Programme of the European Union











Open (Q): Closed: { J (6) } { B (0),

- { B (0), C (2),
- C (2), G (3),

A (4), D (4), Co-funded by the Erasmus+ Programme

of the European Union





J (6) }

Dijkstra's algoritam



Final path solution: $B \rightarrow C \rightarrow D \rightarrow J$ with path cost 6





A* heuristička pretraga

- Heuristika:
 - Funkcija koja opisuje koliko smo daleko od cilja
 - Uvek se postavlja za određeni problem
 - Menhetn ili Euklidska distanca, ...
- A* funkcija cene: f(n) = g(n) + h(n)

Cena puta

Cena heuristike





A* heuristička pretraga







A* heuristička pretraga

Open (Q): {<mark>B(0)</mark>}

Closed: {B(0)}



Our A* queue will be ordered by cost to arrive + heuristic:

- push (Q.Insert) by A* priority, f(n)
- pop (Q.GetFirst) from the front, and add it to the closed list

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5

2

2





Closed:

{ B (0) }

A* heuristička pretraga





Open (Q): { C (2+3), G (3+5), A (5+5) }





A* heuristička pretraga





Open (Q): { G (3+5), A (4+5) } Closed: { B (0) , C (2) }





A* heuristička pretraga





Open (Q): Closed: { D (4+2), { B (0), G (3+5), C (2) } A (4+5) }





A* heuristička pretraga





Open (Q): Closed: { D (4+2), { B (0), G (3+5), C (2) } A (4+5) }





A* heuristička pretraga





Open (Q): Closed: { D (4+2), { B (0) , I (7+1), C (2) } G (3+5), A (4+5) }











ciosca.
{ B (0) ,
C (2),
D (4) }

Closed:

Open (0):





A* heuristička pretraga





Open (Q): Closed: { | (5+1), { B (0), G (3+5), C (2), A (4+5) } D (4) }





Closed:

{ B (0) ,

C (2),

D (4) }



2

R.

Open (0):

{ J (6+0),

I (5+1),

G (3+5),

A (4+5) }



5 2 2 5 5	
A 3 2 1 2 2 2	3





A* heuristička pretraga

2

Closed:

{ B (0),

C (2),

D(4),

J (6) }



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Final path solution: $B \rightarrow C \rightarrow D \rightarrow J$ with path cost 6



https://qiao.github.io/PathFinding.js/visual/









Thanks!

