



Itasdi

Innovative Teaching Approaches in development of Software
Designed Instrumentation and its application in real-time
systems

Theory of Robotics Systems

Planning – getting from A to B

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Innovative Teaching Approaches in development of Software Designed Instrumentation and its application in real-time systems

Faculty of Technical
Sciences



Ss. Cyril and Methodius
University
Faculty of Electrical Engineering
and Information Technologies



Zagreb University of
Applied Sciences



School of Electrical
Engineering
University of Belgrade



Faculty of Physics
Warsaw University of Technology



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Šta znači planiranje?

- Generalno, naš fokus će biti na **planiranju kretanja** u robotici. Određivaćemo set akcija koje će robota odvesti iz poznate pozicije u neku drugu takođe poznatu poziciju.
- Takođe posmatraćemo koja su to **ograničenja** koja mogu uticati na planiranje i od čega potiču kako bismo znali da li je naše planiranje kretanja **izvodljivo**.

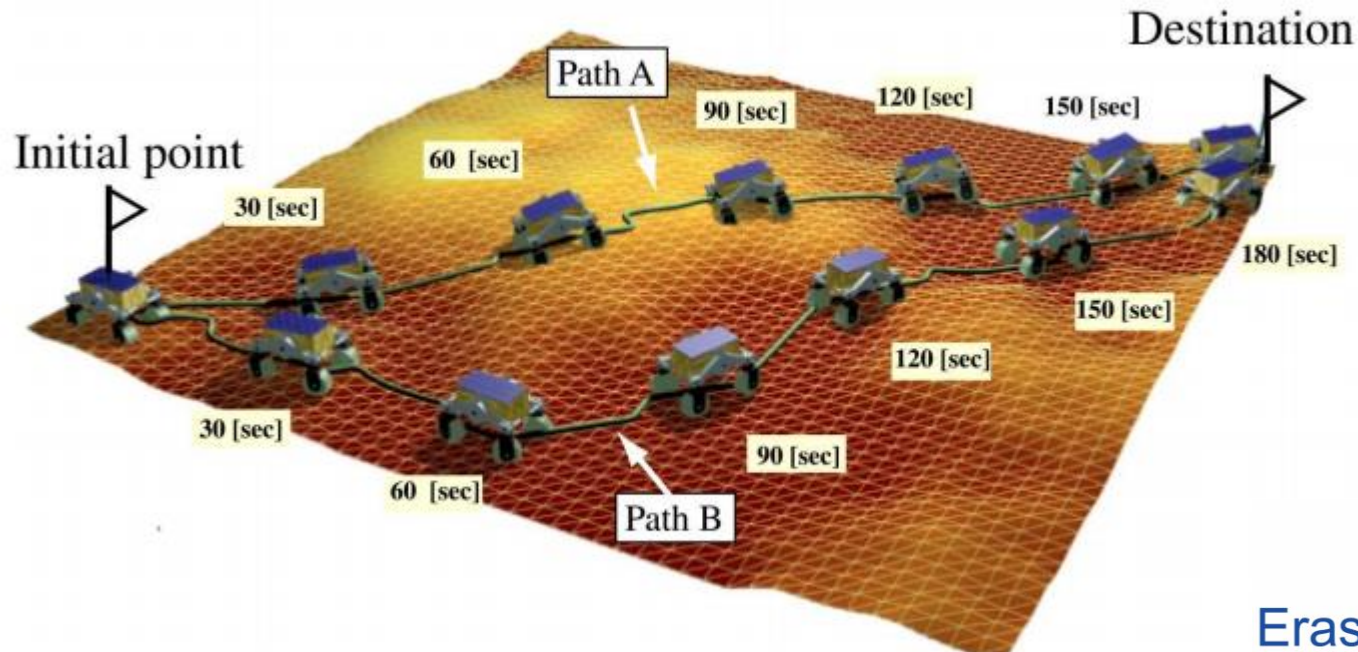




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Šta je planiranje u mobilnoj robotici?



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Povezane teme

- Upravljanje:
 - Generalni zadatak dohvatanja (dolaženja) do željenog stanja.
 - Najčešće koristimo **feedback** upravljanje.
 - Uspešnost upravljanja merimo pomoću stabilnosti, robusnosti i otpornosti na poremećaj.
- Planiranje u Veštačkoj Inteligenciji:
 - Bavi se diskretnim problemima.
 - Klasični algoritmi u VI (pretraga grafova, problem trgovačkog putnika, orijentiring) se često primenjuju kao algoritmi za planiranje u robotici.



FINGER PIVOTING



SLIDING



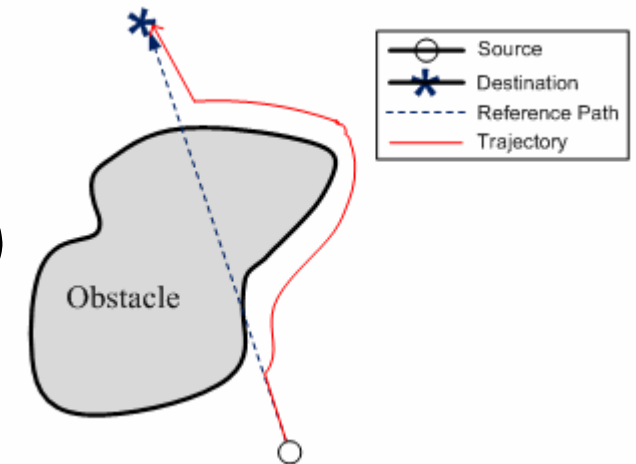
FINGER GAITING





Istorija planiranja kretanja u robotici

- Klasična robotika (sredina 70-ih)
 - Tačan model, bez senzora
- Reaktivno planiranje trajektorije (sredina 80-ih)
 - Bez modela, oslanja se isključivo na sensoriku





Istorija planiranja kretanja u robotici

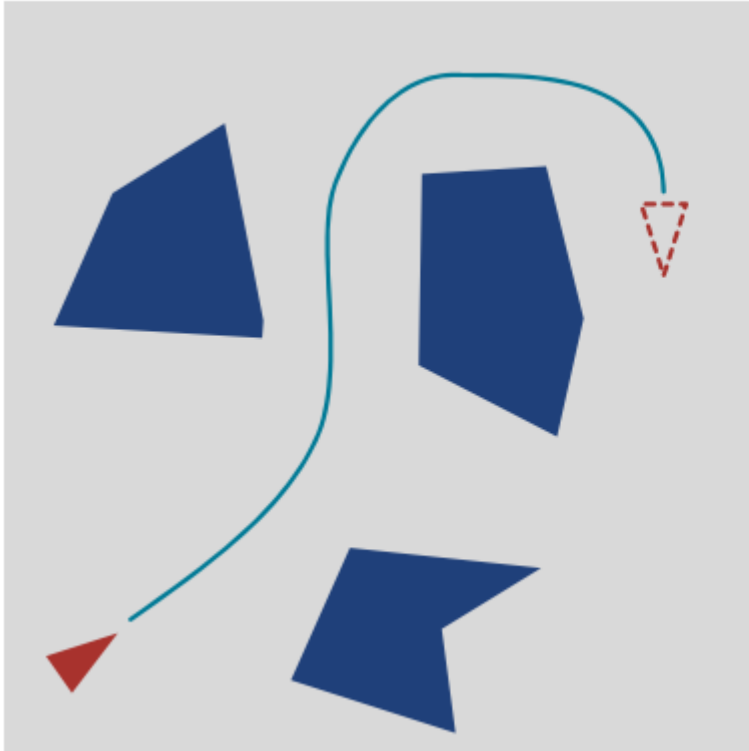
- Hibridno / hijerarhijsko planiranje (početak 90-ih)
 - Model za globalno planiranje, viši nivo
 - Reaktivno (zaobilazanje prepreka), niži nivo
- Probabilističko planiranje (sredina 90-ih)
 - Uključuje nesigurnost modela i senzora u sve nivoi planiranja



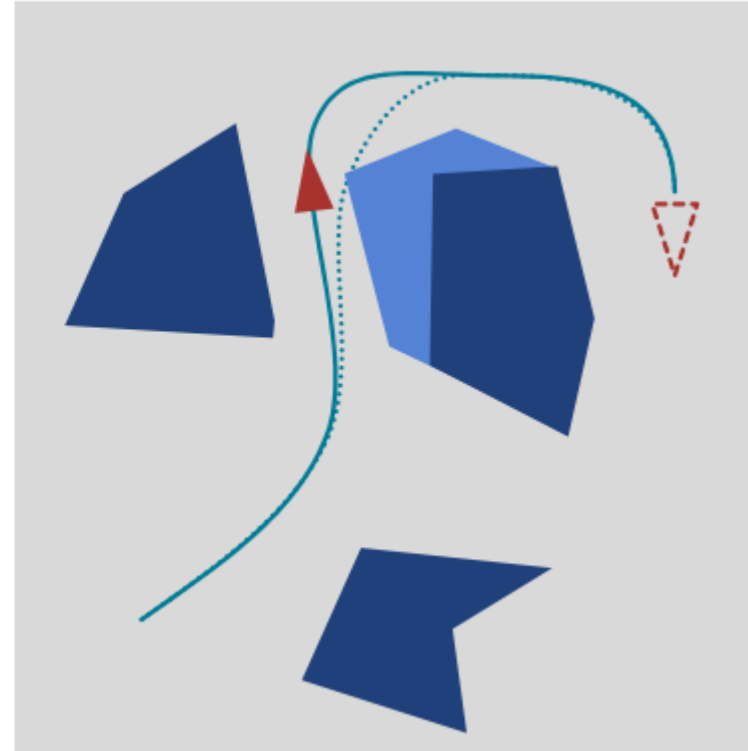


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Planirano



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Uloga navigacije

- Ukratko, naći način kako robot može da se pomeri sa jednog mesta na drugo, poštujući sva ograničenja.
- Neke od pretpostavki:
 - Dovoljno dobro možemo da **predstavimo** okolinu i robota
 - Znamo gde se robot **nalazi** i gde **treba da dođe**
 - Posedujemo **model kretanja** robota
- Obično se predstavlja kao problem optimizacije – minimizacija cene (vreme, udaljenost, energija), poštujući ograničenja.





Planiranje

- Robot se nalazi na startnoj poziciji S , treba da dođe do tačke G :
 - Potrebno je pronaći put koji vodi robota od S do G .
- Izazovi:
 - Prepreke?
 - Dinamika?
 - Izvodljivost?
- Optimization problem !





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Reprezentacija

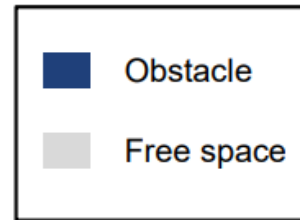
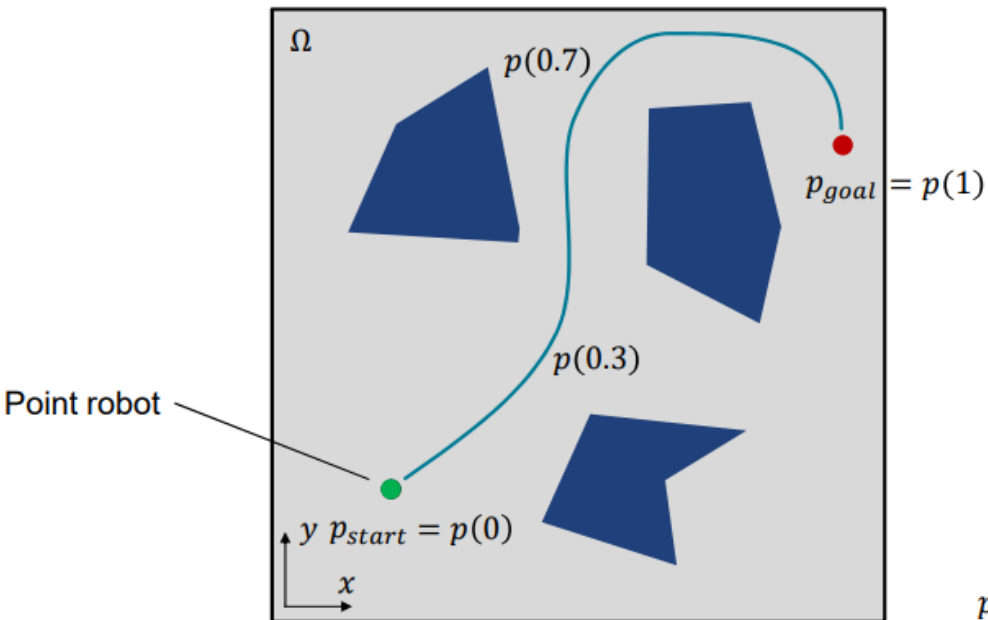
- Veoma je bitno na koji način robot predstavlja i razume svet.
- Uglavnom se moraju primeniti određene metode pojednostavljivanja.
- Odabirom odgovarajuće reprezentacije sveta, određeni algoritmi se mogu primeniti kako bi rešili problem planiranja.

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Reprezentacija – radni prostor i putanja



$$\omega_{free} = \Omega - \omega_{obs}$$

$$p: [0,1] \rightarrow \omega_{free}$$

$$p(t): \forall t \in [0,1], \\ p(t) \in \omega_{free}$$

$$p(0) = p_{start}$$

$$p(1) = p_{goal}$$





Reprezentacija – radni prostor i putanja

- Radni prostor uglavnom predstavlja sam svet (mapu), po mogućstvu nezavisno od robota. Uglavnom na određeni način predstavlja koje su oblasti slobodne a koje ne.
- Konfiguracijski prostor opisuje sva stanja u kojima se robot u datom prostoru može naći.
- Od sad smatramo da naš robot nije tačka u prostoru već sam zauzima određeni prostor.

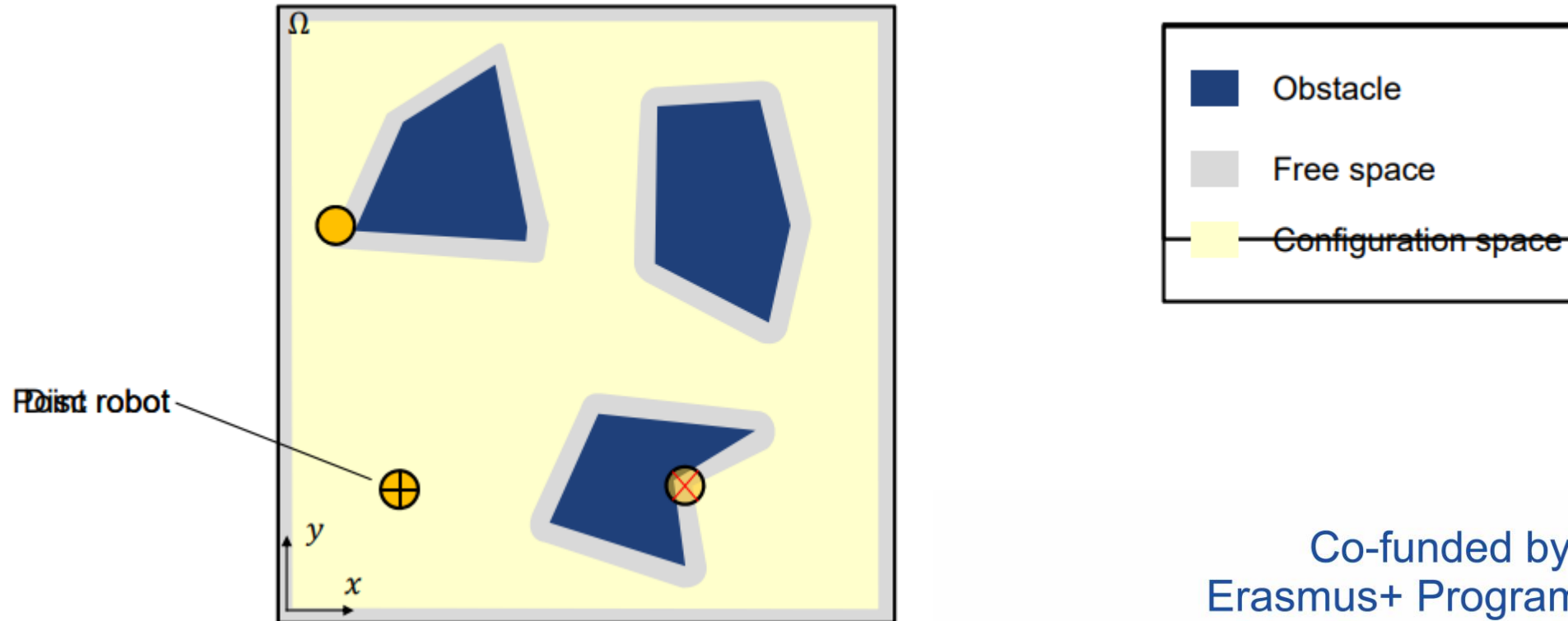




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Konfiguracijski prostor



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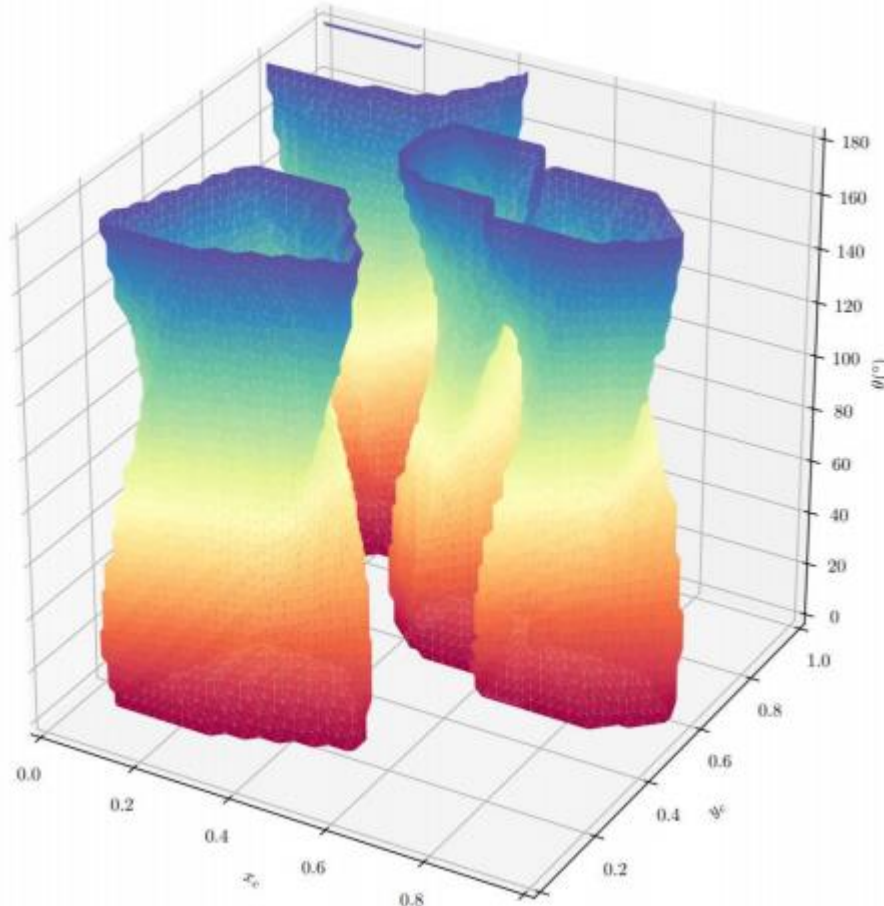




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Konfiguracijski prostor



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Kontinualna vs Diskretna reprezentacija

- Iako za kontinualnu reprezentaciju posedujemo dobar matematički alat, uglavnom nam takva predstava nije zgodna.
- Kako računari čuvaju podatke digitalno, potrebno je naći pametan i efikasan način za kreiranje (aproksimacije) diskretne reprezentacije.
- Veoma je uobičajeno da problem planiranja predstavimo kao problem pretrage grafa, jer za pretragu grafova postoji dosta već razrađenih algoritama.

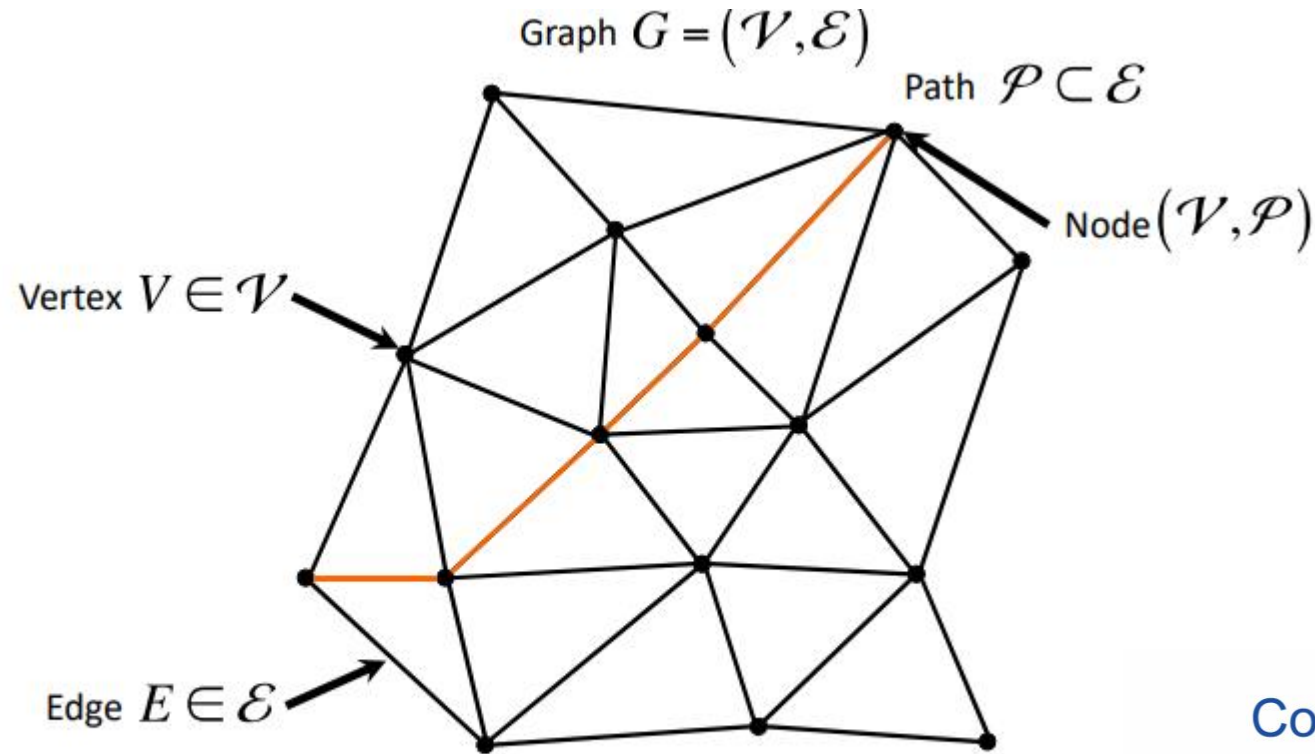




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Terminologija



Directed graph: edges have direction
Weighted graph: edges have costs

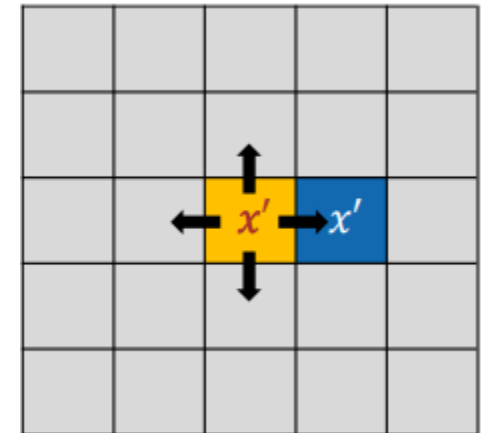
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Diskretna reprezentacija prostora stanja

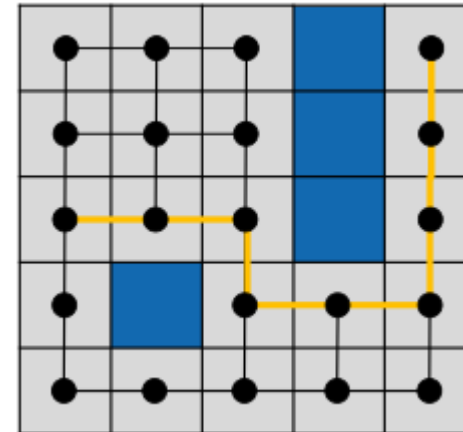
- Redukujemo kontinualna stanja na skup konačnih diskretnih stanja.
 - $x \in X$
- Definišemo izvodljiv skup akcija za svako stanje
 - $A(x) = \{a_0, a_1, \dots, a_n\}$
- I svakoj od akcija pridružujemo funkciju tranzicije
 - $f(x, a) = x'$





Grid \rightarrow Graf

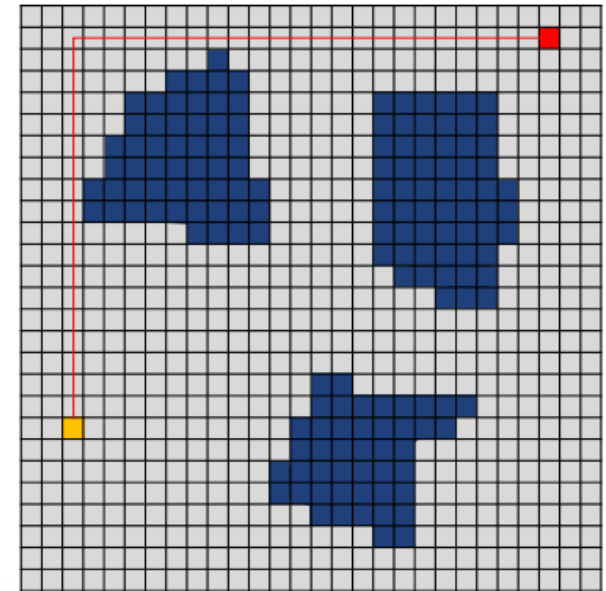
- Primeniti:
 - Stanja se uzimaju kao vrhovi
 - Prelazi kao usmerene ivice
- Rezultat je graf
- Dodajemo:
 - Početni nod, x_s
 - Završni nod, x_g
 - Funkcija cene $C: X \times A \rightarrow \mathbb{R}^+$
- Traženje najkraćeg puta možemo tretirati kao problem pretrage grafa.





Problemi sa reprezentacijom pomoću grid-a

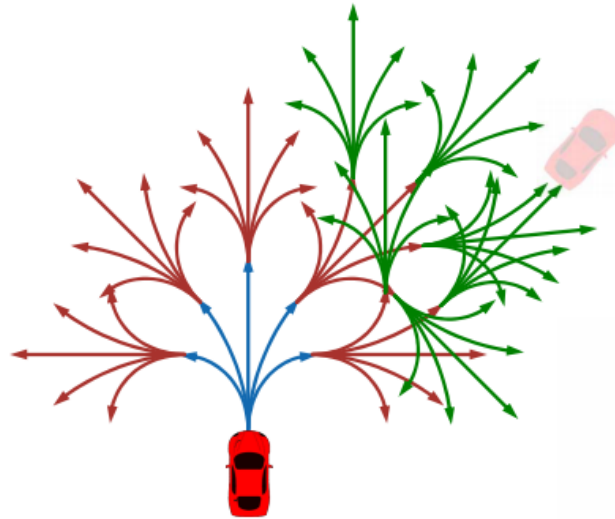
- Uglavnom se neki sitni detalji i preciznost mogu izgubiti.
- Odabir prave rezolucije može da bude veoma izazovan zadatak (mapiranje multi-rezolucijom)
- Izlazna putanja može biti ograničena.
- Loša rezolucija za mape velikih dimenzija.





Grid rešetka

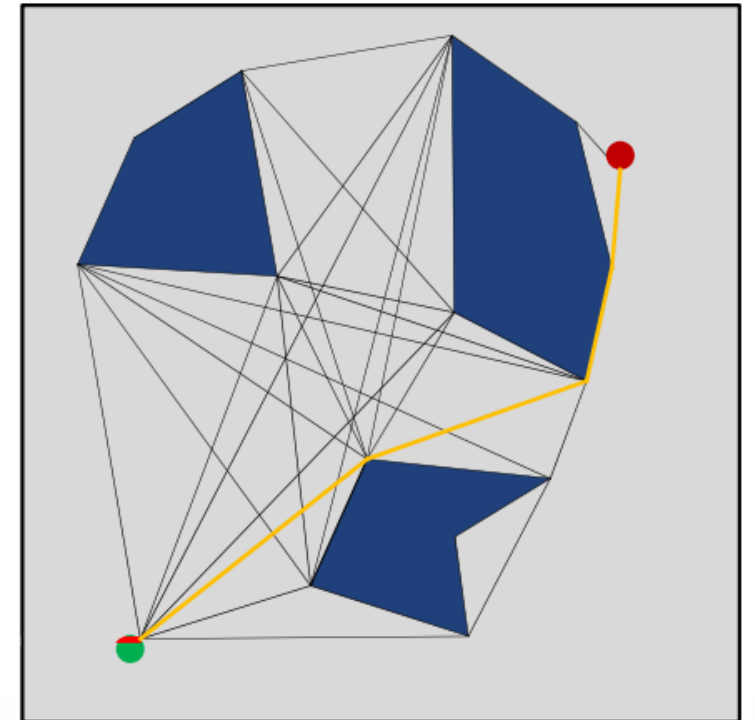
- Kreira set mogućih primitivnih kretanja, i konstruiše drvo (graf) koje sadrži određena kretanja u sekvenci (putanja)





Grafikon vidljivosti

- Kreira ivice grafa između svih parova vidljivih čvorova.
- Pretraga se vrši prema dobijenom grafu.
- Optimalan put je putanja kretanja.
- Ograničeno kretanje na prave putanje.



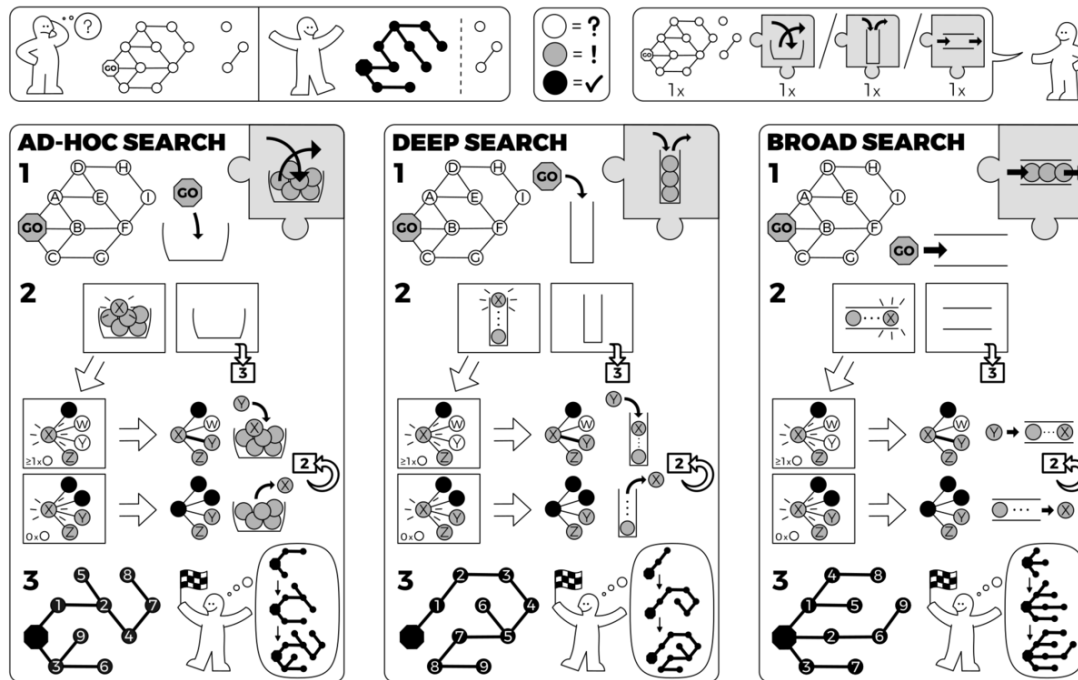


Planiranje kretanja

- Planiranje kretanja: deljenje zadatka kretanja na diskretna kretanja:
 - Robot mora da stigne od početne lokacija na cilj.
 - Robot mora da zaobiđe zidove, da uspešno izbegne stepenice, ...
 - Idealno, putanja mora da bude najefikasnija (najkraća)
- Ulazi algoritma:
 - Opis zadatka (početna i krajnja pozicija)
 - Ograničenja (fizička ograničenja robota, prepreke, ...)
 - Nesigurnosti robota i okruženja
- Izlaz algoritma:
 - Putanja robota
 - Set tačaka koje treba da poseti kako bi otišao od tačke A do tačke B



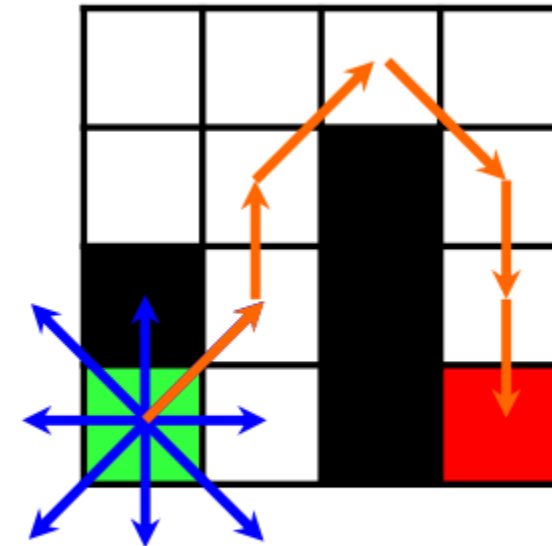
Metode pretrage grafova





Primer 1

- Putanja od zelene ćelije do crvene
- 8 koraka iz jedne ćelije
- Koja je najkraća putanja?
- Šta ako imamo 4 moguća koraka iz ćelije.



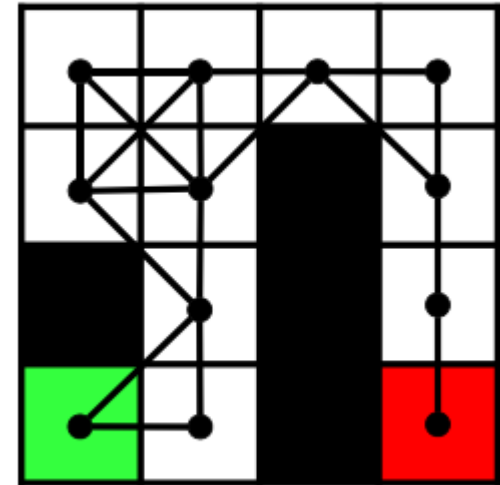
Assumptions about
robot motion





Primer 1

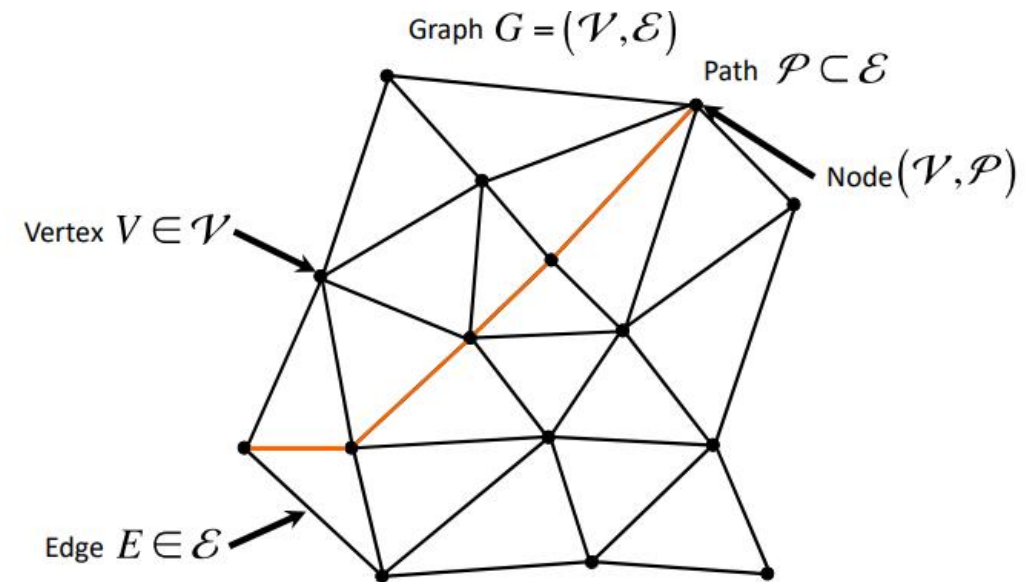
- Pretražimo dobijeni graf.
- Pretpostavljamo da sve ivice grafa mogu kreirati putanju.
- Rešimo problem tako što nađemo putanju od A do B.





Pretraga grafova

- Ivice grafa mogu imati pridružene **težine**
 - Cena putanje
- Ima smisla da se razmatraju samo pozitivne težine.
- Algoritmi sa minimalnom cenom nikad ne posećuju isti čvor dva puta.
- Najpoznatije pretrage:
 - Pretraga u širinu (BFS)
 - Pretraga u dubinu (DFS)
 - Dijkstra's Algoritam
 - A*



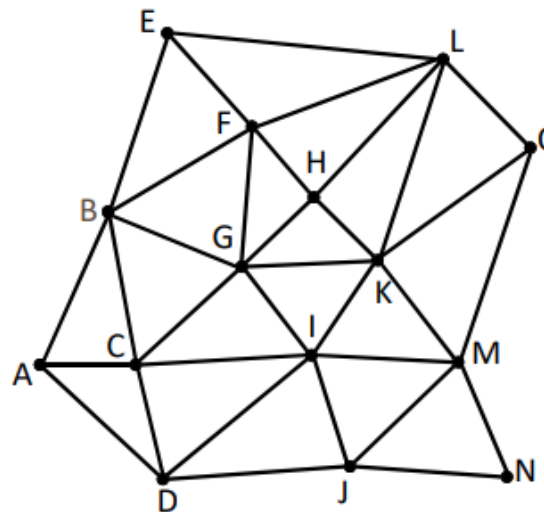
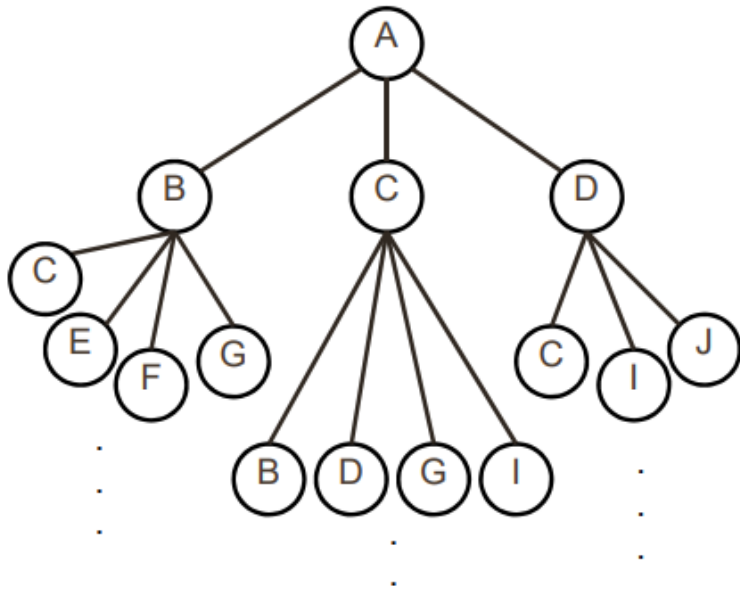
Directed graph: edges have direction
Weighted graph: edges have costs





Drvo pretrage

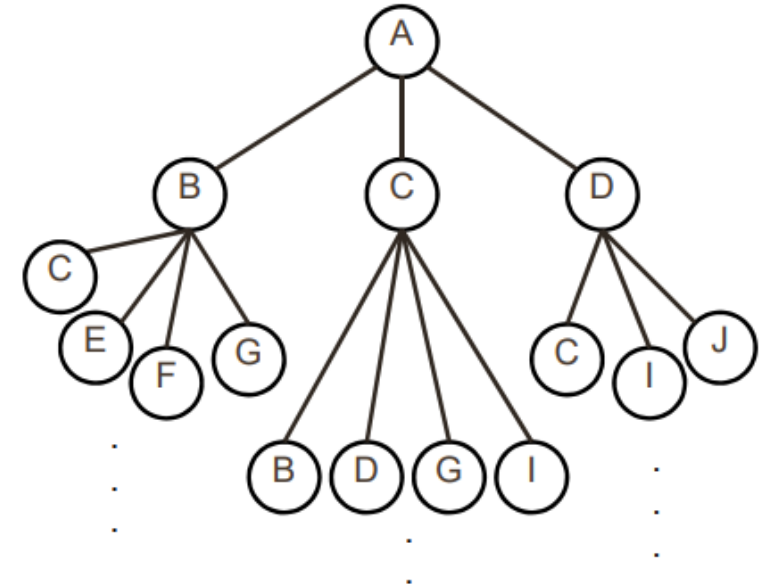
- Konstruišemo „drvo“ pomoću kojeg vršimo pretragu.





Drvo pretrage

- Drvo pretrage:
 - Početno stanje se uzima kao koren
 - Naslednici – skup svih mogućih akcija
 - Uglavnom želimo da izbegnemo konstrukciju kompletnog stabla





Osnove pravolinijske pretrage

- Počnemo od korena (početno stanje), širimo drvo sve dok ne dođemo do cilja.
- Proširivanje nodova znači dodavanje njihovih naslednika.
- Treba probati da se nađe rešenje sa što manje proširivanja.
- Otvoren set – čvorovi iz kojih se mogu vršiti proširivanja
- Zatvoren set – svi oni čvorovi koje je algoritam posetio





Osnove pravolinijske pretrage

```
FORWARD_SEARCH
1  Q.Insert( $x_I$ ) and mark  $x_I$  as visited
2  while  $Q$  not empty do
3       $x \leftarrow Q.GetFirst()$ 
4      if  $x \in X_G$ 
5          return SUCCESS
6      forall  $u \in U(x)$ 
7           $x' \leftarrow f(x, u)$ 
8          if  $x'$  not visited
9              Mark  $x'$  as visited
10             Q.Insert( $x'$ )
11         else
12             Resolve duplicate  $x'$ 
13 return FAILURE
```

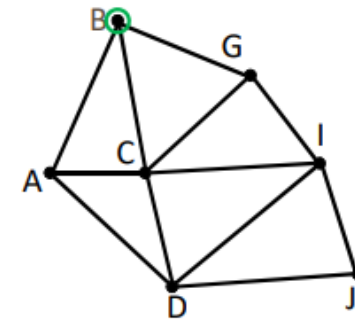




Pretraga po širini

FORWARD SEARCH

```
1 Q.Insert( $x_1$ ) and mark  $x_1$  as visited
2 while Q not empty do
3    $x \leftarrow Q.GetFirst()$ 
4   if  $x \in X_G$ 
5     return SUCCESS
6   forall  $u \in U(x)$ 
7      $x' \leftarrow f(x, u)$ 
8     if  $x'$  not visited
9       Mark  $x'$  as visited
10      Q.Insert( $x'$ )
11   else
12     Resolve duplicate  $x'$ 
13 return FAILURE
```



Open (Q): Closed:
{B} {}

Our (BFS) queue will be FIFO:

- push (*Q.Insert*) onto the end
- pop (*Q.GetFirst*) from the front



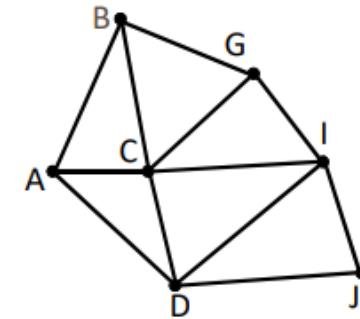
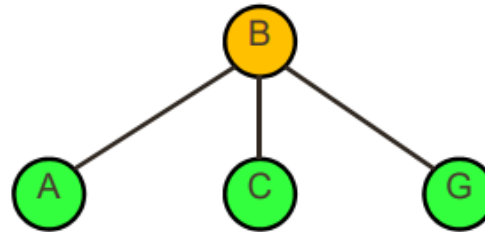


Pretraga po širini

```
FORWARD_SEARCH
1  Q.Insert(xI) and mark xI as visited
2  while Q not empty do
3    x ← Q.GetFirst()
4    if x ∈ XG
5      return SUCCESS
6    forall u ∈ U(x)
7      x' ← f(x, u)
8      if x' not visited
9        Mark x' as visited
10       Q.Insert(x')
11     else
12       Resolve duplicate x'
13  return FAILURE
```

Open (Q):
{A,C,G}

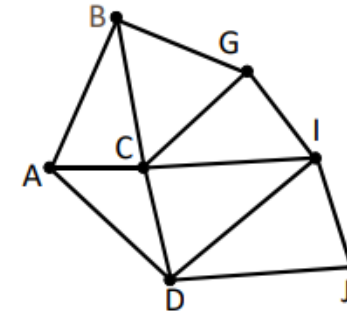
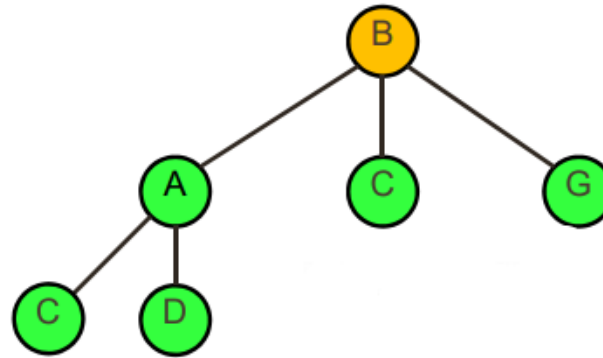
Closed:
{B,A,C,G}





Pretraga po širini

```
FORWARD_SEARCH
1  Q.Insert(xl) and mark xl as visited
2  while Q not empty do
3    x ← Q.GetFirst()
4    if x ∈ XG
5      return SUCCESS
6    forall u ∈ U(x)
7      x' ← f(x, u)
8      if x' not visited
9        Mark x' as visited
10       Q.Insert(x')
11     else
12       Resolve duplicate x'
13  return FAILURE
```



Open (Q):

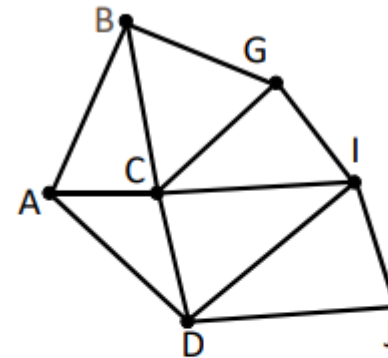
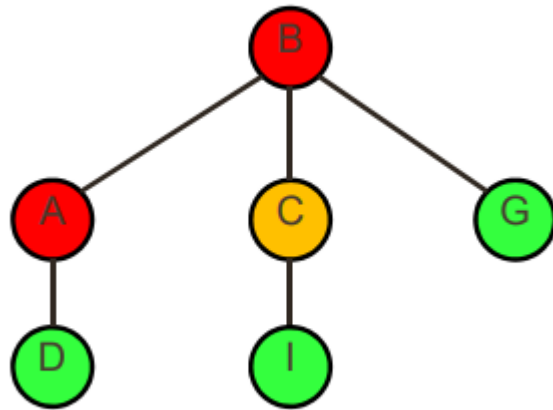
{C,G,D}

Closed:

{B,A,C,G,D}



Pretraga po širini

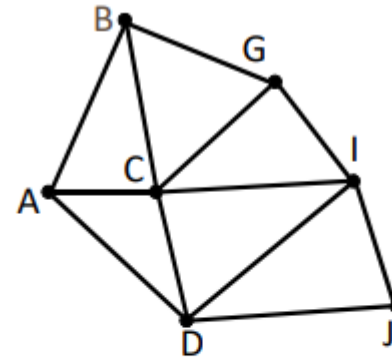
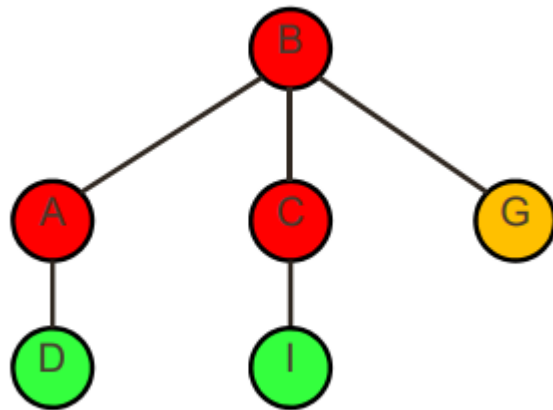


Open (Q):
{G,D,I}

Closed:
{B,A,C,G,D,I}



Pretraga po širini

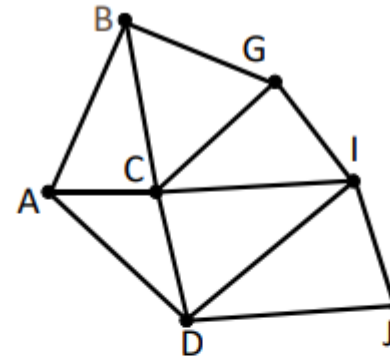
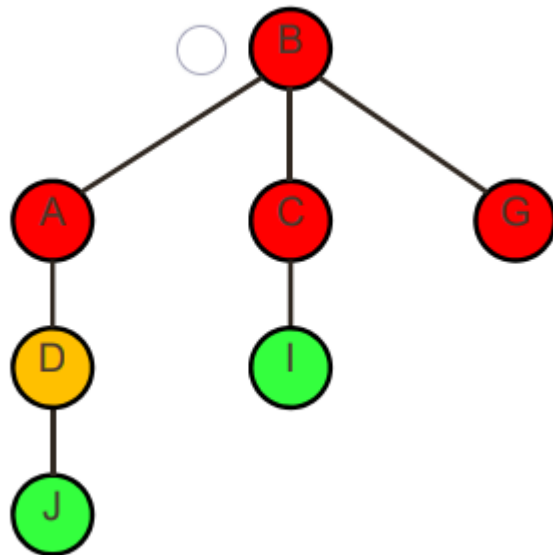


Open (Q):
{D,I}

Closed:
{B,A,C,G,D,I}



Pretraga po širini

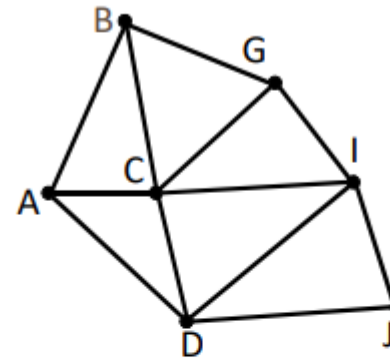
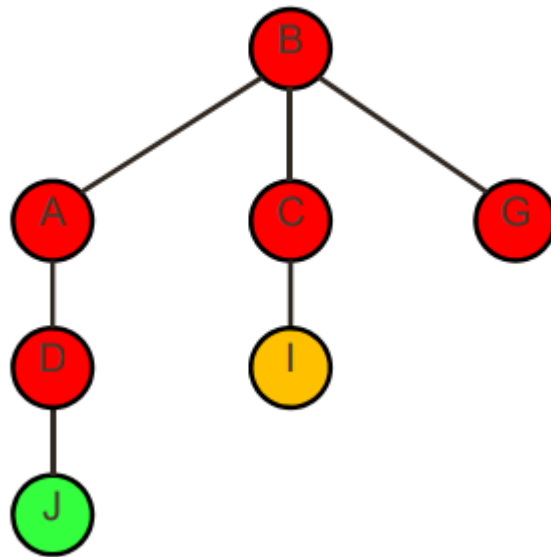


Open (Q):
{I,J}

Closed:
{B,A,C,G,D,I,J}



Pretraga po širini

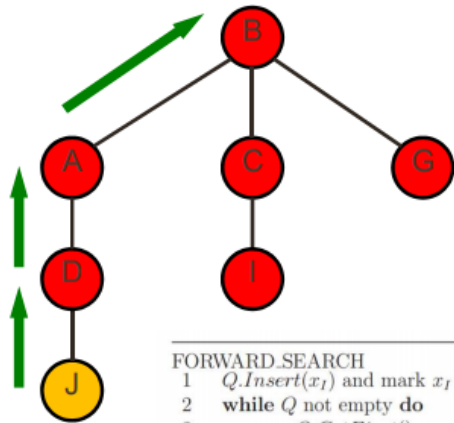


Open (Q):
{J}

Closed:
{B,A,C,G,D,I,J}

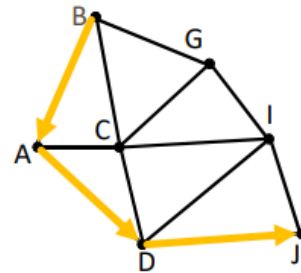


Pretraga po širini



```

FORWARD_SEARCH
1  Q.Insert(xi) and mark xi as visited
2  while Q not empty do
3    x ← Q.GetFirst()
4    if x ∈ XG
5      return SUCCESS
6    forall u ∈ U(x)
7      x' ← f(x, u)
8      if x' not visited
9        Mark x' as visited
10       Q.Insert(x')
11     else
12       Resolve duplicate x'
13  return FAILURE
  
```



Open (Q): {}
 Closed: {B,A,C,G,D,I,J}

Final path solution: B → A → D → J

Other solutions may exist but have the same number or more transitions





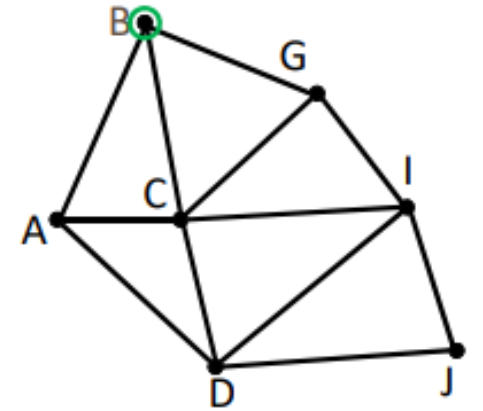
Pretraga po dubini

FORWARD SEARCH

```

1  Q.Insert( $x_f$ ) and mark  $x_f$  as visited
2  while Q not empty do
3     $x \leftarrow Q.GetFirst()$ 
4    if  $x \in X_G$ 
5      return SUCCESS
6    forall  $u \in U(x)$ 
7       $x' \leftarrow f(x, u)$ 
8      if  $x'$  not visited
9        Mark  $x'$  as visited
10     Q.Insert( $x'$ )
11   else
12     Resolve duplicate  $x'$ 
13 return FAILURE

```



Open (Q):

{B}

Closed:

{}

- Our (DFS) queue will be LIFO:
- push (*Q.Insert*) onto the front
 - pop (*Q.GetFirst*) from the front

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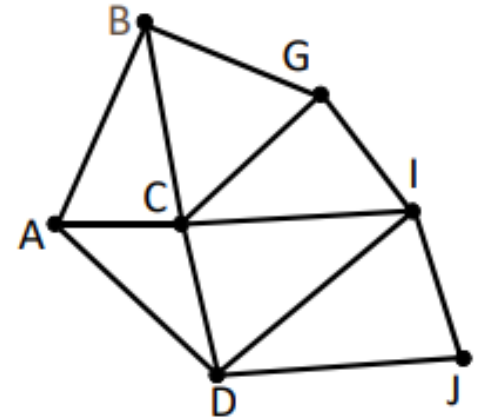
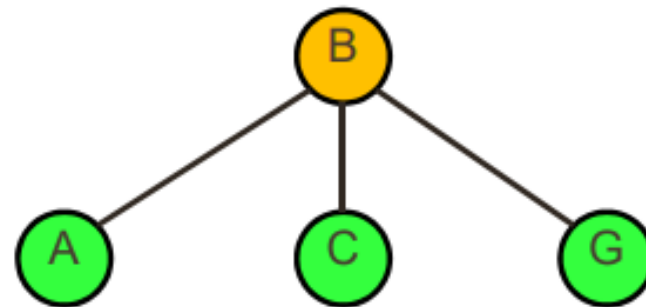




Pretraga po dubini

FORWARD_SEARCH

```
1  Q.Insert( $x_I$ ) and mark  $x_I$  as visited
2  while Q not empty do
3     $x \leftarrow Q.GetFirst()$ 
4    if  $x \in X_G$ 
5      return SUCCESS
6    forall  $u \in U(x)$ 
7       $x' \leftarrow f(x, u)$ 
8      if  $x'$  not visited
9        Mark  $x'$  as visited
10       Q.Insert( $x'$ )
11     else
12       Resolve duplicate  $x'$ 
13  return FAILURE
```



Open (Q):
{A,C,G}

Closed:
{B,A,C,G}

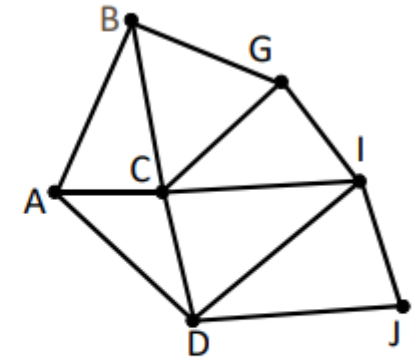
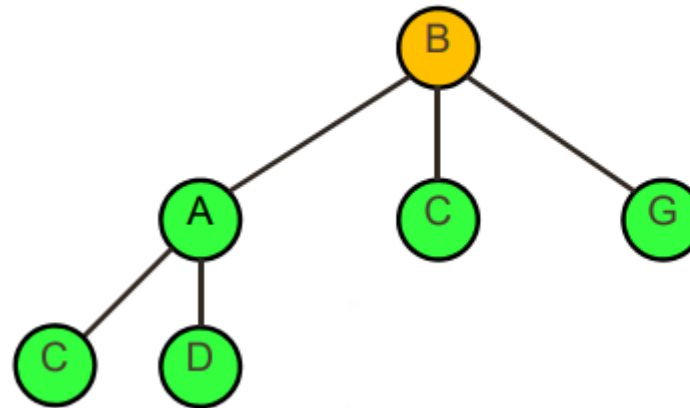




Pretraga po dubini

FORWARD_SEARCH

```
1  Q.Insert( $x_I$ ) and mark  $x_I$  as visited
2  while Q not empty do
3     $x \leftarrow Q.GetFirst()$ 
4    if  $x \in X_G$ 
5      return SUCCESS
6    forall  $u \in U(x)$ 
7       $x' \leftarrow f(x, u)$ 
8      if  $x'$  not visited
9        Mark  $x'$  as visited
10       Q.Insert( $x'$ )
11     else
12       Resolve duplicate  $x'$ 
13  return FAILURE
```



Open (Q):
{D,C,G}

Closed:
{B,A,C,G,D}

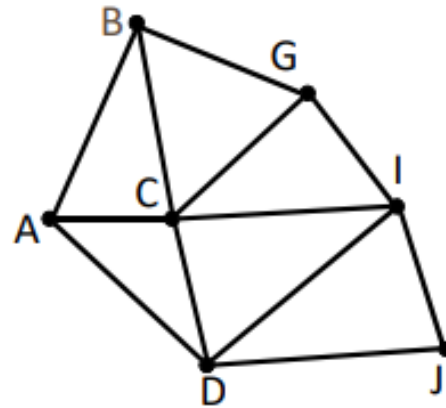
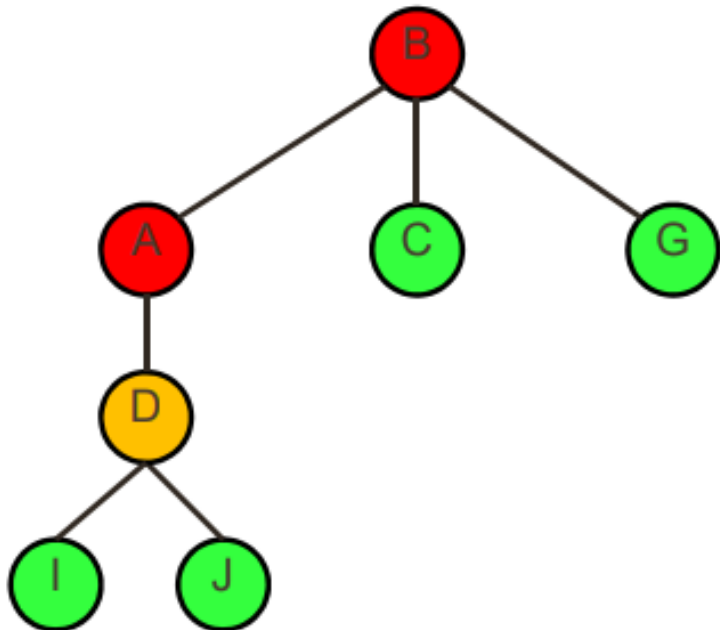




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Pretraga po dubini



Open (Q):
{I,J,C,G}

Closed:
{B,A,C,G,D,I,J}

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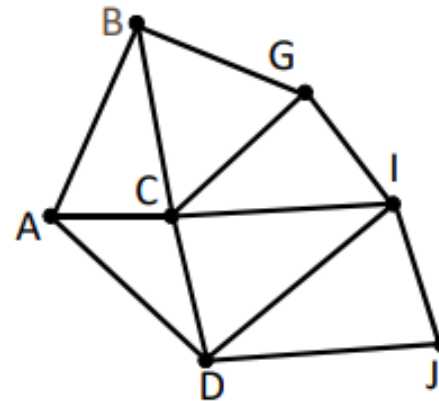
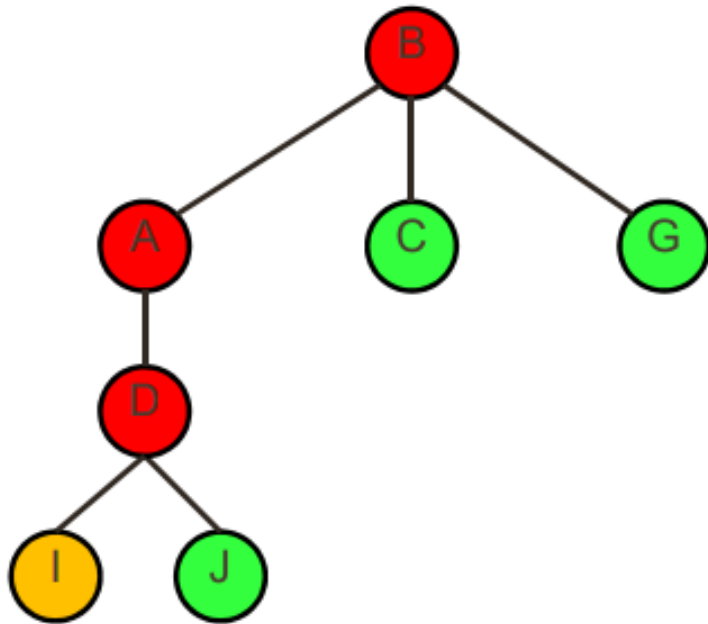




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Pretraga po dubini



Open (Q):
{J,C,G}

Closed:
{B,A,C,G,D,I,J}

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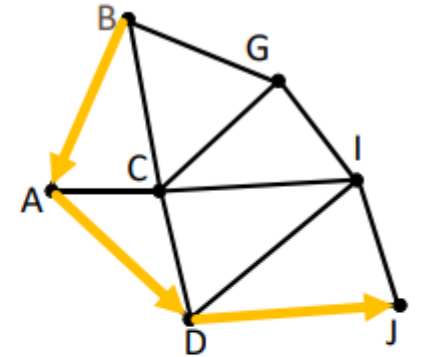
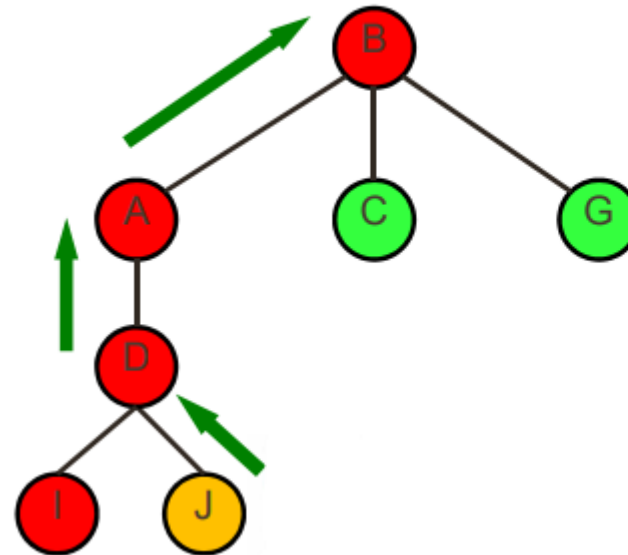
Pretraga po dubini

FORWARD-SEARCH

```
1  Q.Insert( $x_I$ ) and mark  $x_I$  as visited
2  while Q not empty do
3     $x \leftarrow Q.GetFirst()$ 
4    if  $x \in X_G$ 
5      return SUCCESS
6    forall  $u \in U(x)$ 
7       $x' \leftarrow f(x, u)$ 
8      if  $x'$  not visited
9        Mark  $x'$  as visited
10     Q.Insert( $x'$ )
11   else
12     Resolve duplicate  $x'$ 
13  return FAILURE
```

Open (Q):
{}

Closed:
{B,A,C,G,D,I,J}



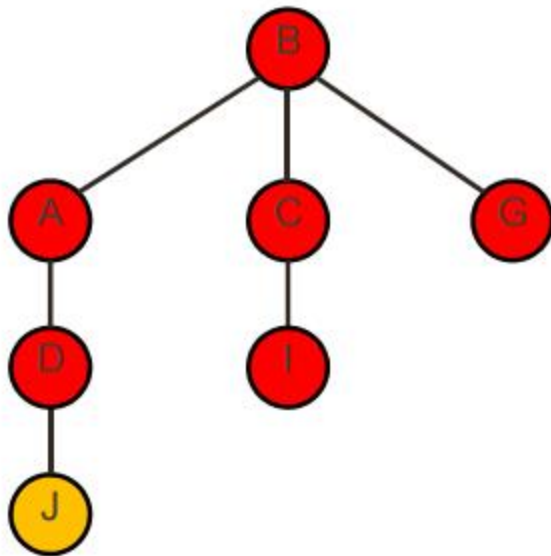
Final path solution: B→A→D→J

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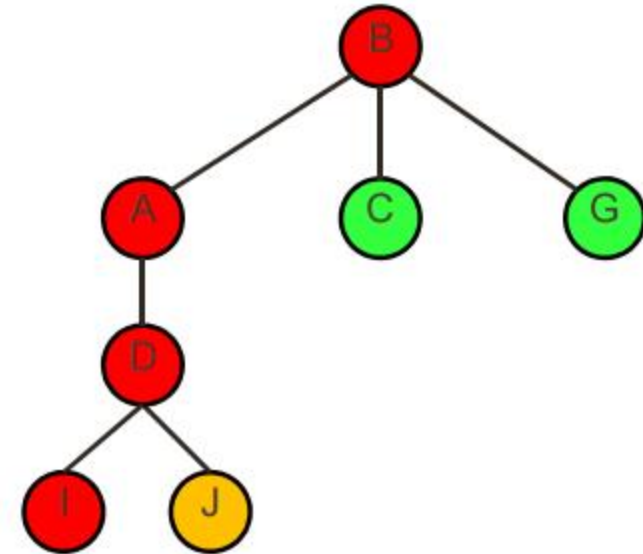
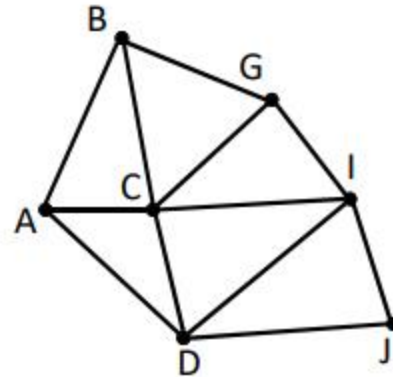




Uporedna analiza



BFS



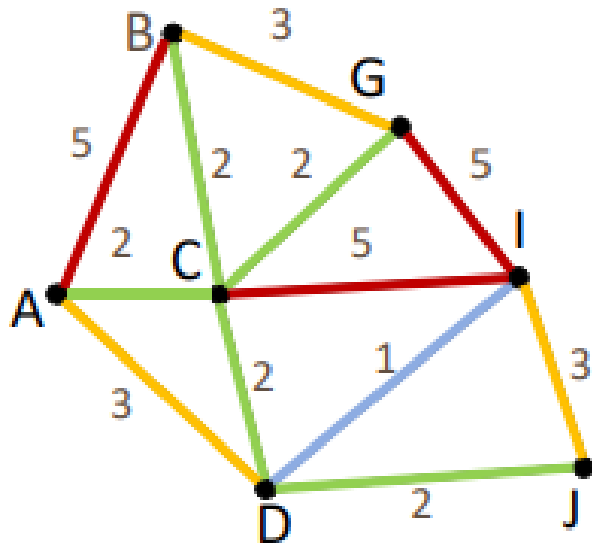
DFS





Akcije sa cenom

- Kako rešiti slučaj ako određena putanja ima veću cenu nego neka druga?
- Najpopularniji algoritmi Dijkstra's i A^*





Dijkstra's algoritam

- Objavljen od strane Edsger Dijsktra 1959.
- Otvara čvorove sa najmanjom cenom najbliže od korena grafa.
- Najkorišćeniji algoritam za rešavanje problema trgovačkog putnika.





Dijkstra's algoritam

FORWARD SEARCH

```
1 Q.Insert( $x_I$ ) and mark  $x_I$  as visited
2 while  $Q$  not empty do
3    $x \leftarrow Q.GetFirst()$ 
4   if  $x \in X_G$ 
5     return SUCCESS
6   forall  $u \in U(x)$ 
7      $x' \leftarrow f(x, u)$ 
8     if  $x'$  not visited
9       Mark  $x'$  as visited
10      Q.Insert( $x'$ )
11   else
12     Resolve duplicate  $x'$ 
13 return FAILURE
```

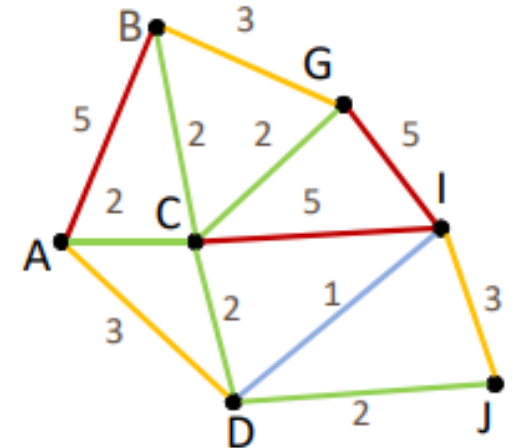
B(0)

Open (Q):
{B(0)}

Closed:
{B(0)}

Our Dijkstra queue will be ordered by cost to arrive:

- push (*Q.Insert*) by cost
- pop (*Q.GetFirst*) from the front, and add it to the closed list





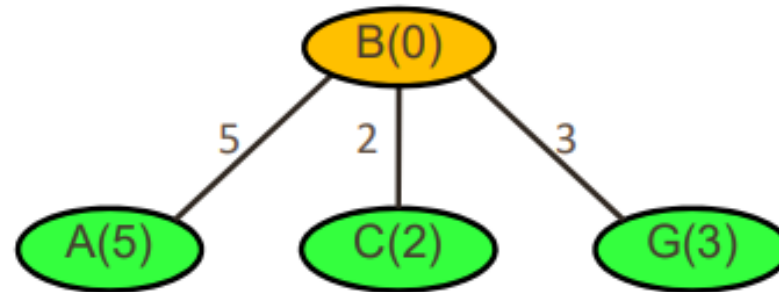
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Dijkstra's algoritam

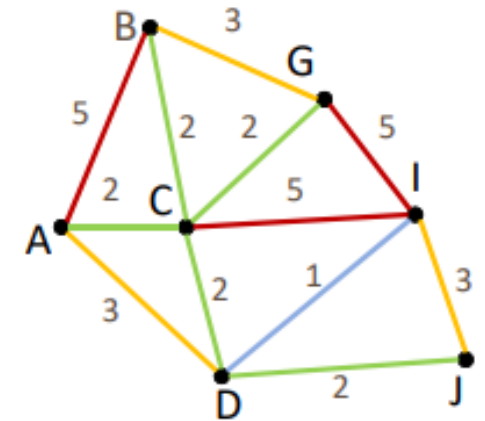
FORWARD_SEARCH

```
1  Q.Insert(xf) and mark xf as visited
2  while Q not empty do
3    x ← Q.GetFirst()
4    if x ∈ XG
5      return SUCCESS
6    forall u ∈ U(x)
7      x' ← f(x, u)
8      if x' not visited
9        Mark x' as visited
10     Q.Insert(x')
11   else
12     Resolve duplicate x'
13  return FAILURE
```



Open (Q):
{ C (2),
G (3),
A (5) }

Closed:
{ B (0) }



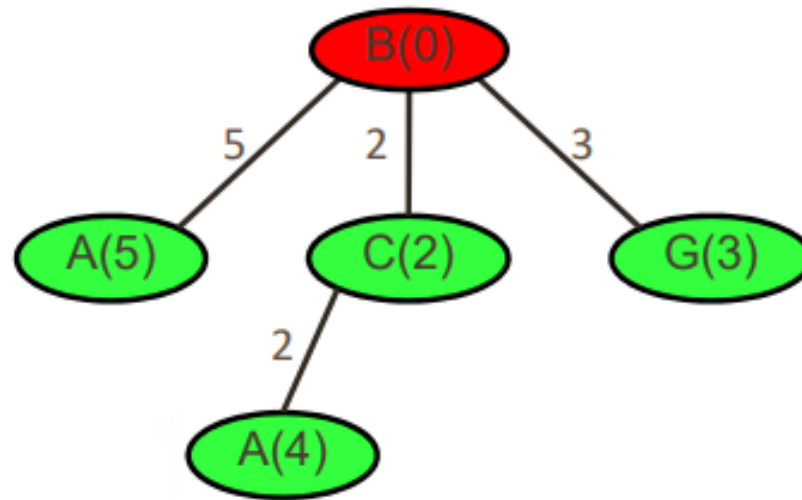
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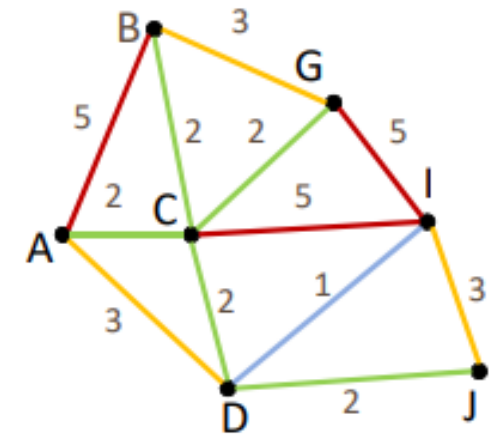
Dijkstra's algoritam

```
FORWARD_SEARCH
1  Q.Insert(xf) and mark xf as visited
2  while Q not empty do
3    x ← Q.GetFirst()
4    if x ∈ XG
5      return SUCCESS
6    forall u ∈ U(x)
7      x' ← f(x,u)
8      if x' not visited
9        Mark x' as visited
10     Q.Insert(x')
11   else
12     Resolve duplicate x'
13  return FAILURE
```



Open (Q):
{ G (3),
A (4) }

Closed:
{ B (0),
C (2) }



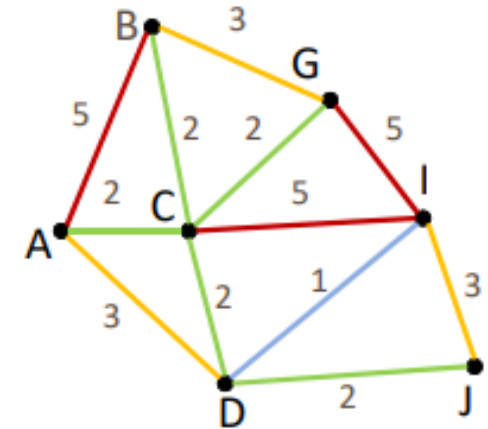
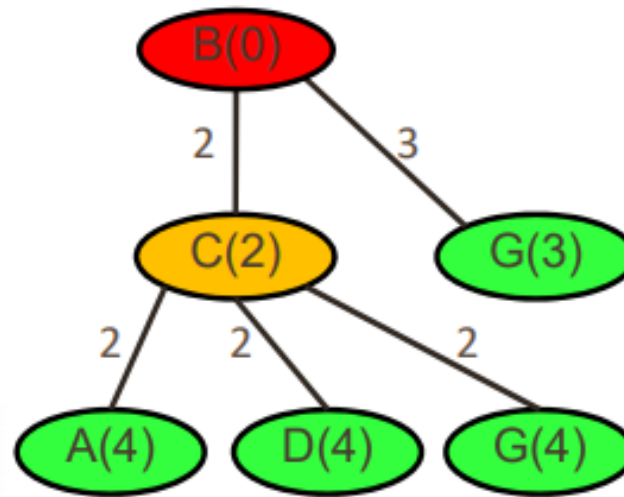


Dijkstra's algoritam

```

FORWARD_SEARCH
1  Q.Insert(xI) and mark xI as visited
2  while Q not empty do
3    x ← Q.GetFirst()
4    if x ∈ XG
5      return SUCCESS
6    forall u ∈ U(x)
7      x' ← f(x,u)
8      if x' not visited
9        Mark x' as visited
10       Q.Insert(x')
11    else
12      Resolve duplicate x'
13  return FAILURE

```



Open (Q):
{ G (3),
A (4),
D (4) }

Closed:
{ B (0),
C (2) }

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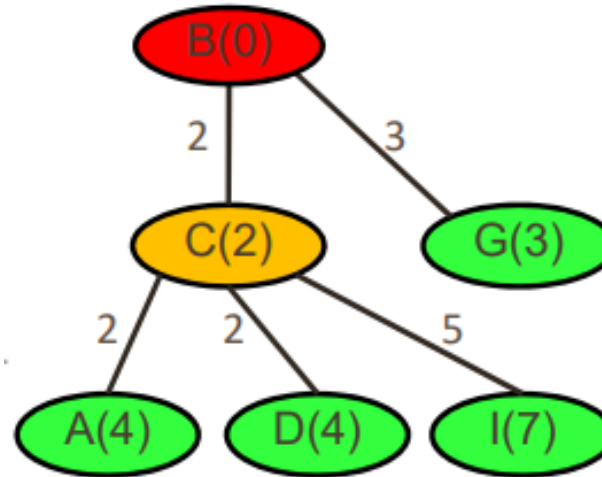
Dijkstra's algoritam

FORWARD_SEARCH

```

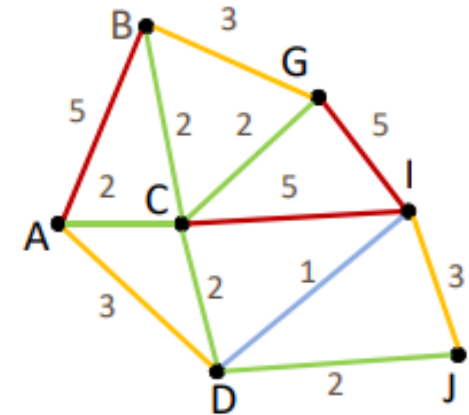
1  Q.Insert(xI) and mark xI as visited
2  while Q not empty do
3    x ← Q.GetFirst()
4    if x ∈ XG
5      return SUCCESS
6    forall u ∈ U(x)
7      x' ← f(x, u)
8      if x' not visited
9        Mark x' as visited
10       Q.Insert(x')
11    else
12      Resolve duplicate x'
13  return FAILURE

```



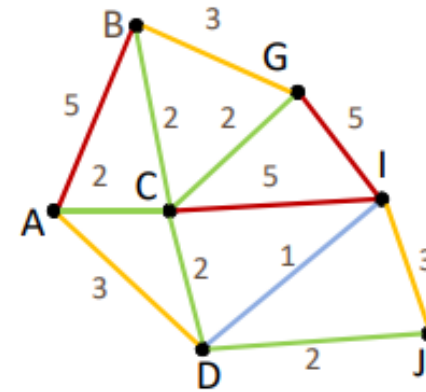
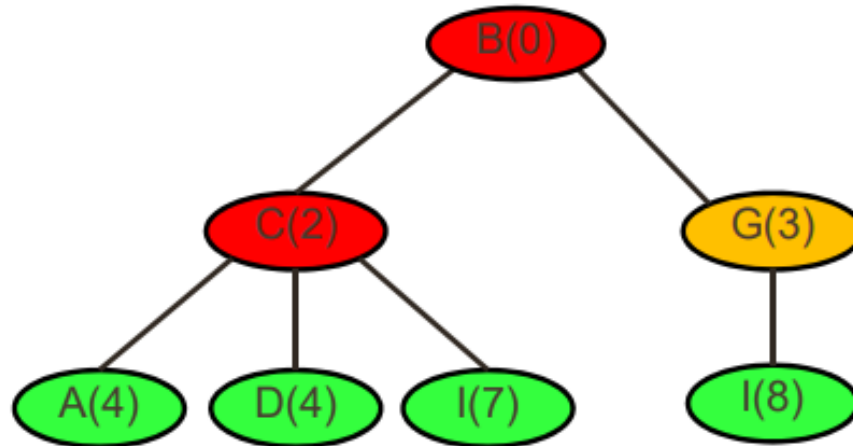
Open (Q):
 { G (3),
 A (4),
 D (4),
 I (7) }

Closed:
 { B (0),
 C (2) }





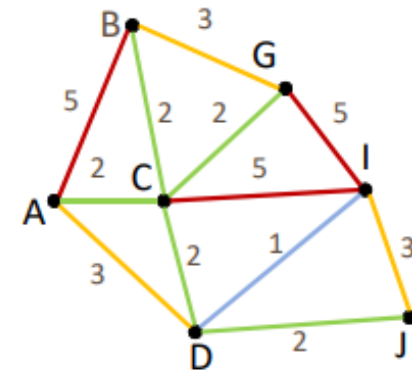
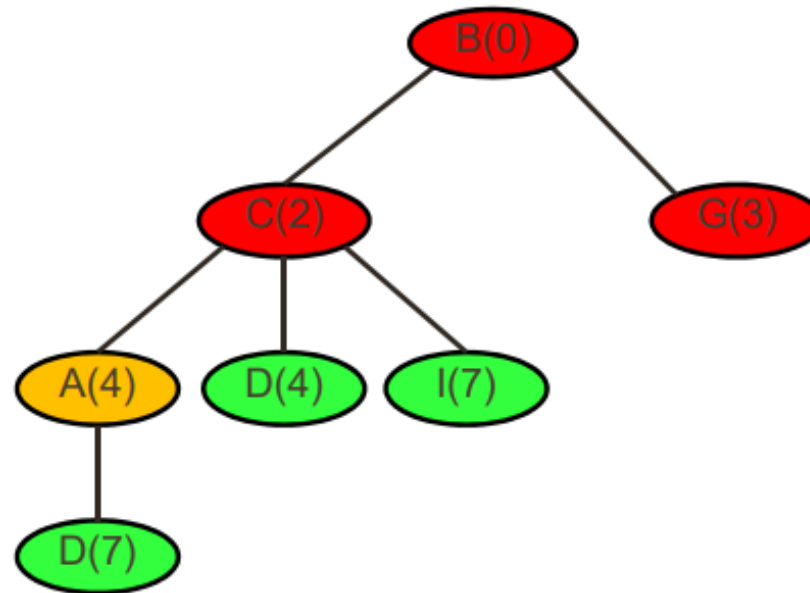
Dijkstra's algoritam



Open (Q):	Closed:
{ A (4),	{ B (0),
D (4),	C (2),
I (7) }	G (3) }



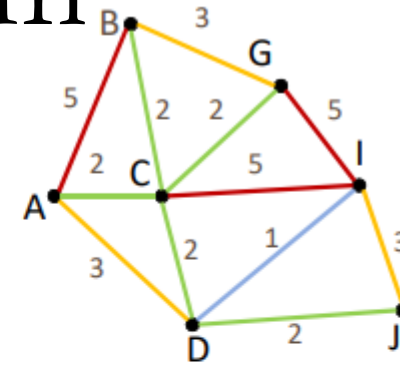
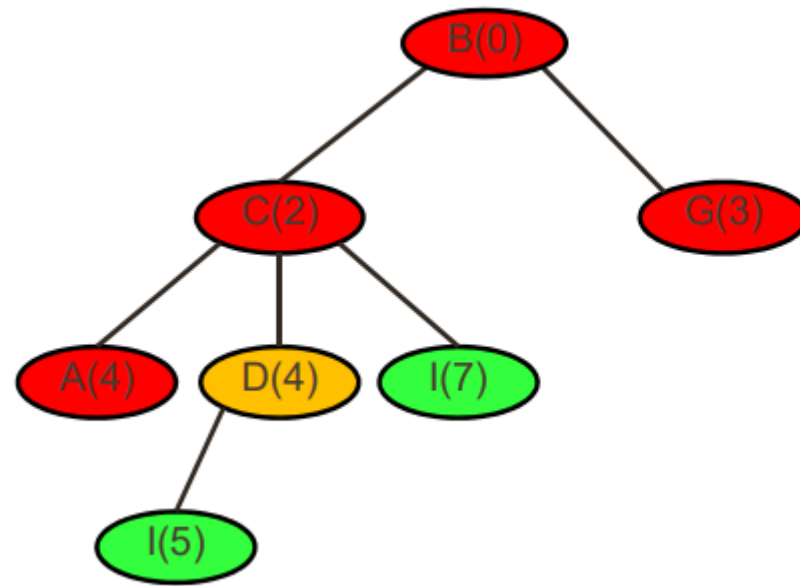
Dijkstra's algoritam



Open (Q):	Closed:
{ D (4),	{ B (0),
I (7) }	C (2),
	G (3),
	A (4) }



Dijkstra's algoritam



Open (Q):	Closed:
{ I (5) }	{ B (0), C (2), G (3), A (4), D (4) }

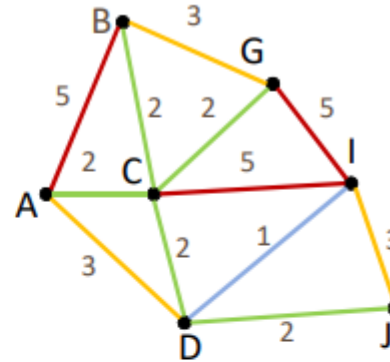
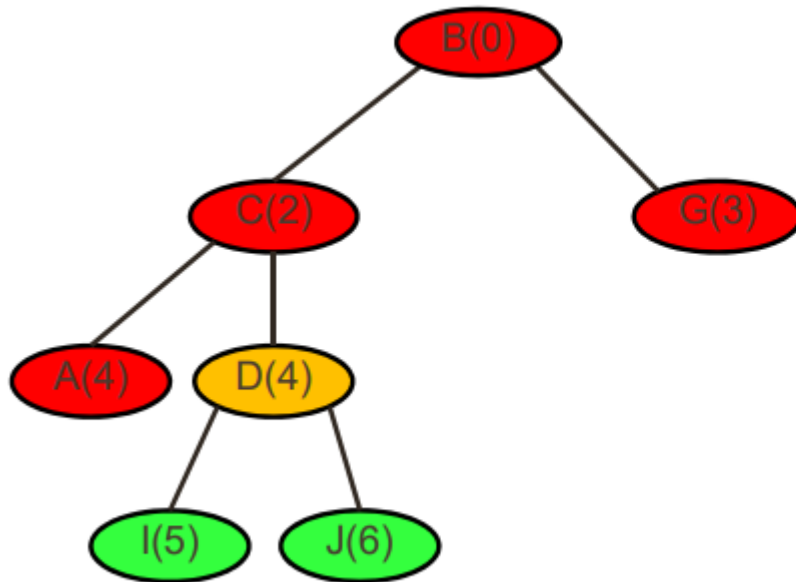




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Dijkstra's algoritam



Open (Q):
{ I (5),
J (6) }

Closed:
{ B (0),
C (2),
G (3),
A (4),
D (4) }

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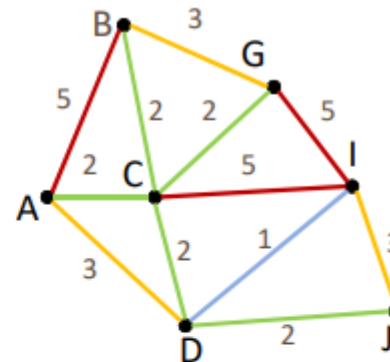
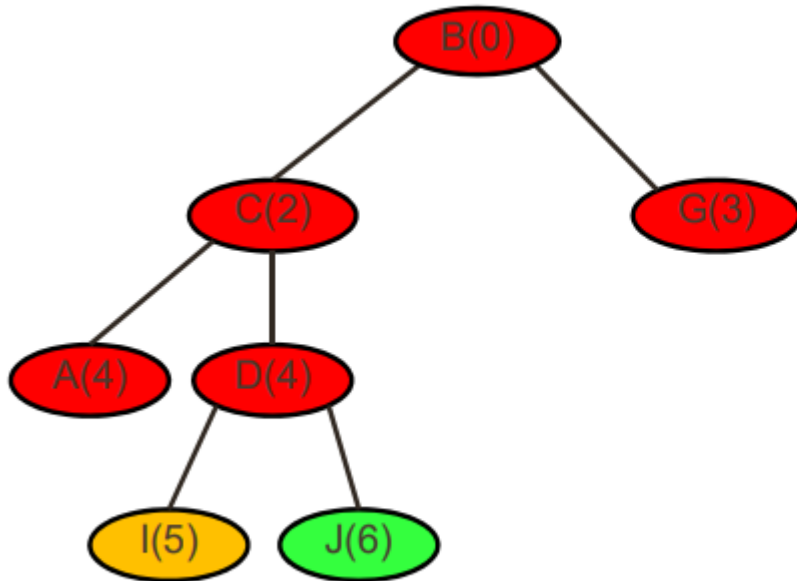




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Dijkstra's algoritam



Open (Q):
{ J (6) }

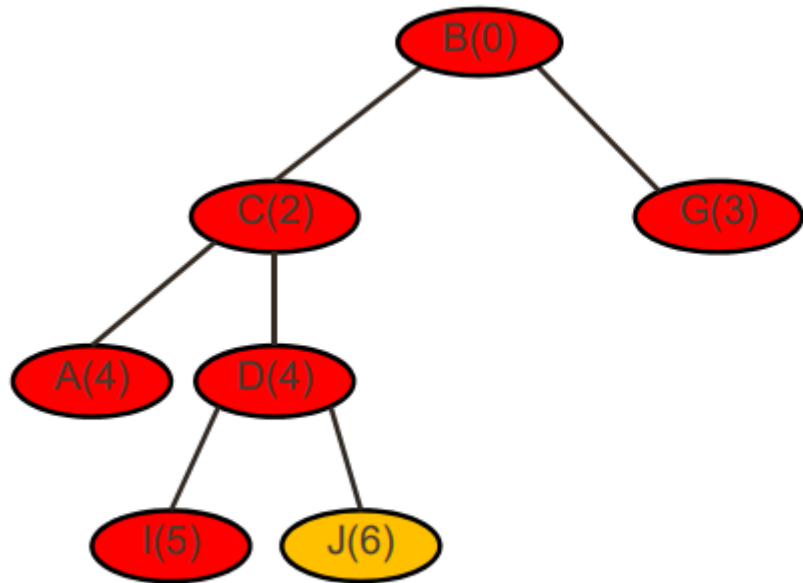
Closed:
{ B (0),
C (2),
G (3),
A (4),
D (4),
I (5) }

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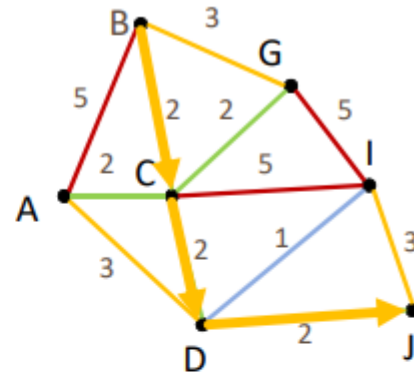




Dijkstra's algoritam



Final path solution: B → C → D → J
with path cost 6



Open (Q):
{ }

Closed:
{ B (0),
C (2),
G (3),
A (4),
D (4),
I (5),
J (6) }





A* heuristička pretraga

- Heuristika:
 - Funkcija koja opisuje koliko smo daleko od cilja
 - Uvek se postavlja za određeni problem
 - Menhetn ili Euklidska distanca, ...
- A* funkcija cene: $f(n) = g(n) + h(n)$

Cena puta

Cena heuristike

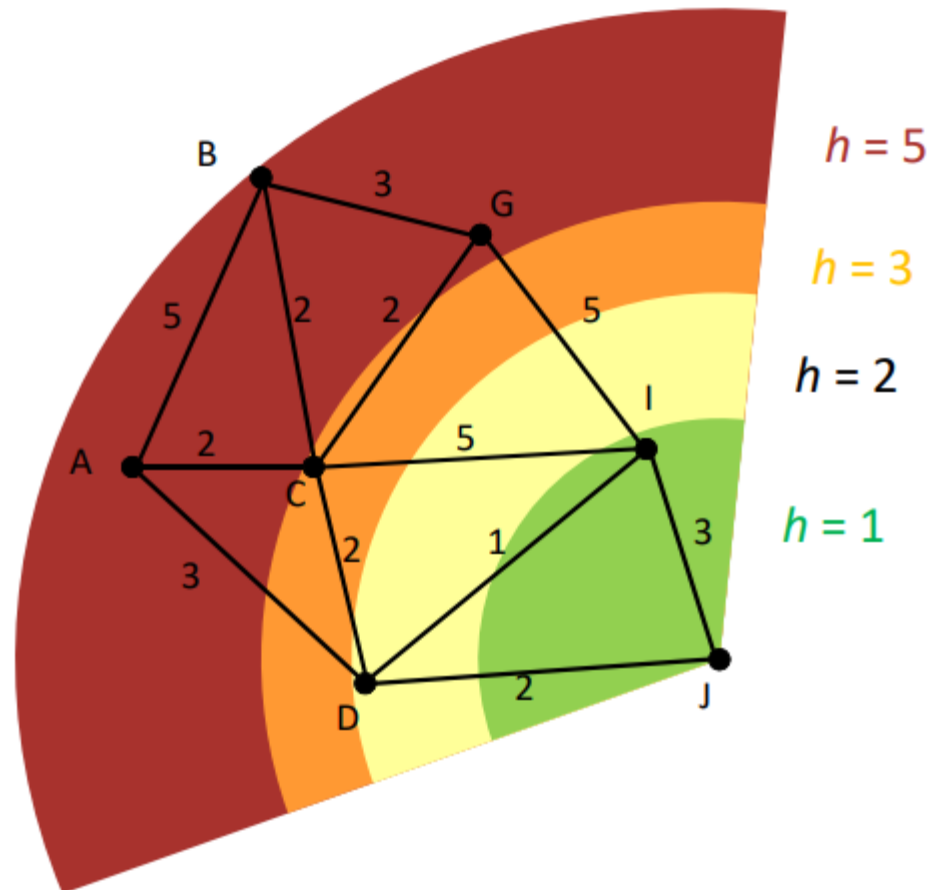




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A* heuristička pretraga



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A* heuristička pretraga

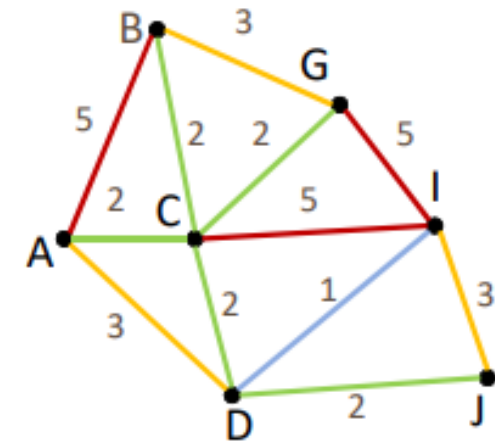
Open (Q):
{B(0)}

Closed:
{B(0)}

B(0)

Our A* queue will be ordered by cost to arrive + heuristic:

- push (*Q.Insert*) by A* priority, $f(n)$
- pop (*Q.GetFirst*) from the front, and add it to the closed list

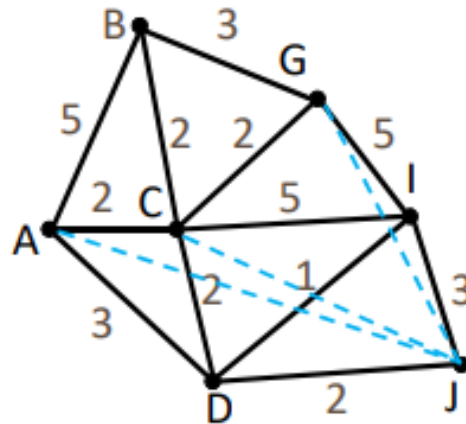
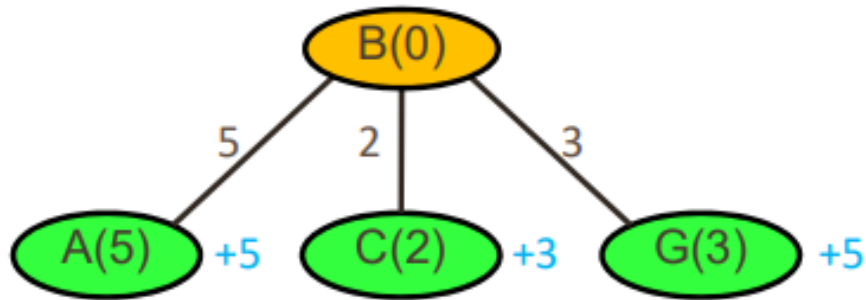


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A* heuristička pretraga



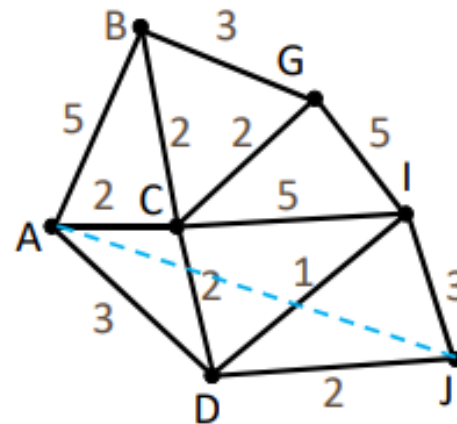
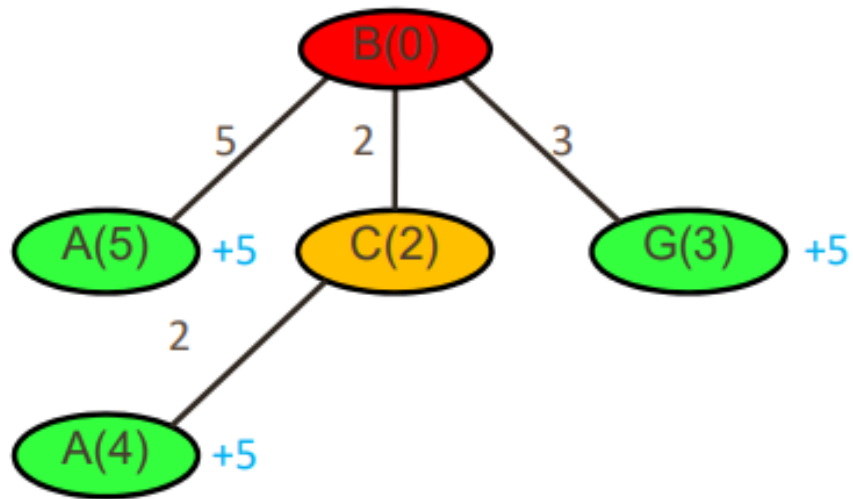
Open (Q):
{ C (2+3),
G (3+5),
A (5+5) }

Closed:
{ B (0) }





A* heuristička pretraga



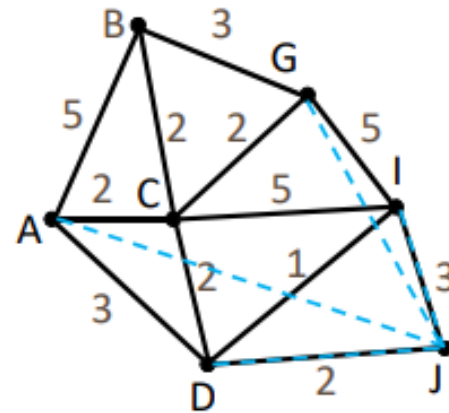
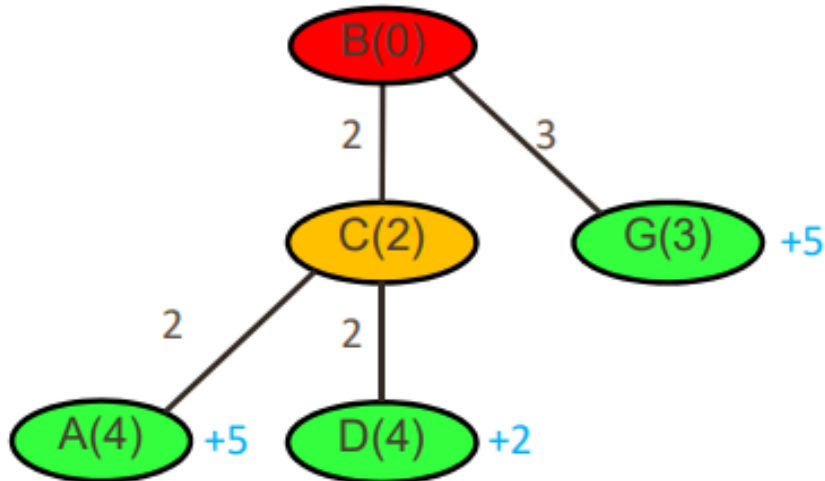
Open (Q):
{ G (3+5),
A (4+5) }

Closed:
{ B (0),
C (2) }





A* heuristička pretraga



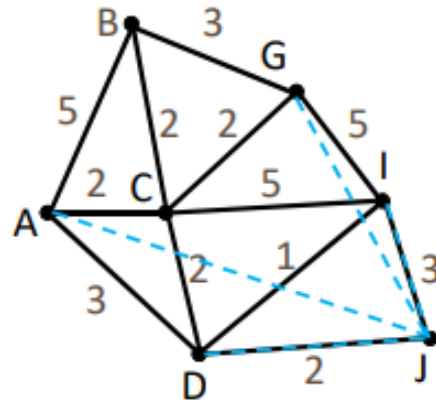
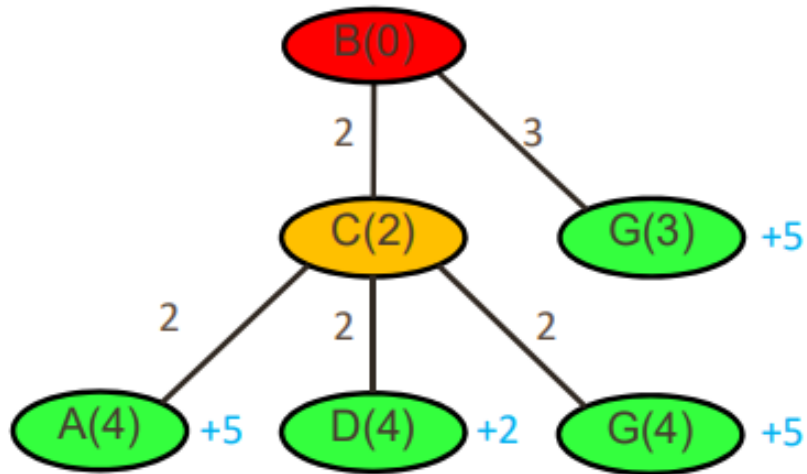
Open (Q):
{ D (4+2),
G (3+5),
A (4+5) }

Closed:
{ B (0),
C (2) }





A* heuristička pretraga



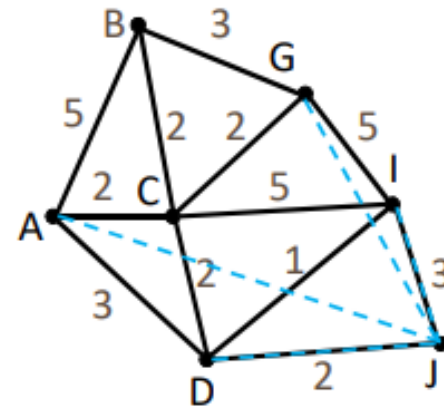
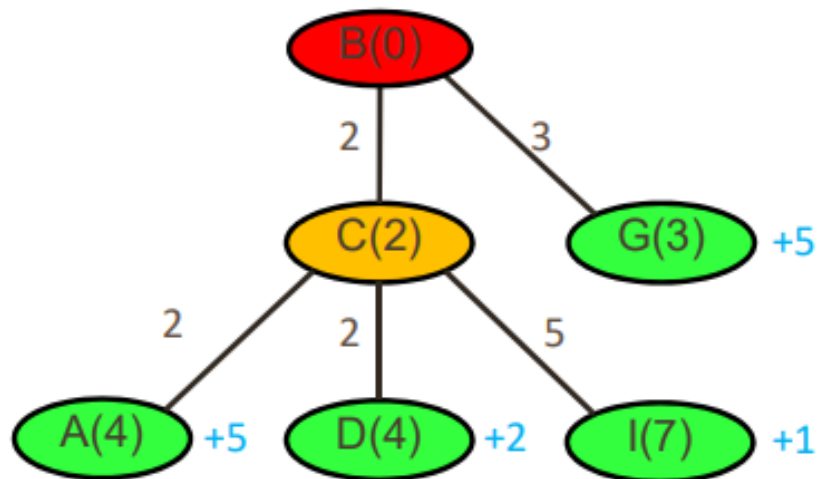
Open (Q):
{ D (4+2),
G (3+5),
A (4+5) }

Closed:
{ B (0),
C (2) }





A* heuristička pretraga



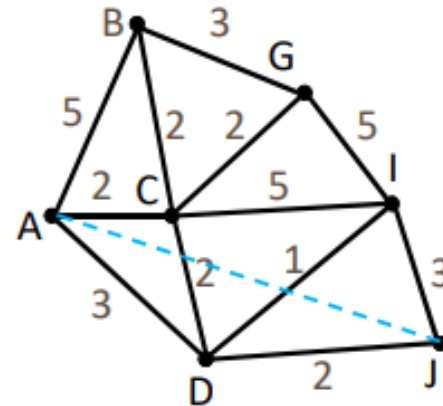
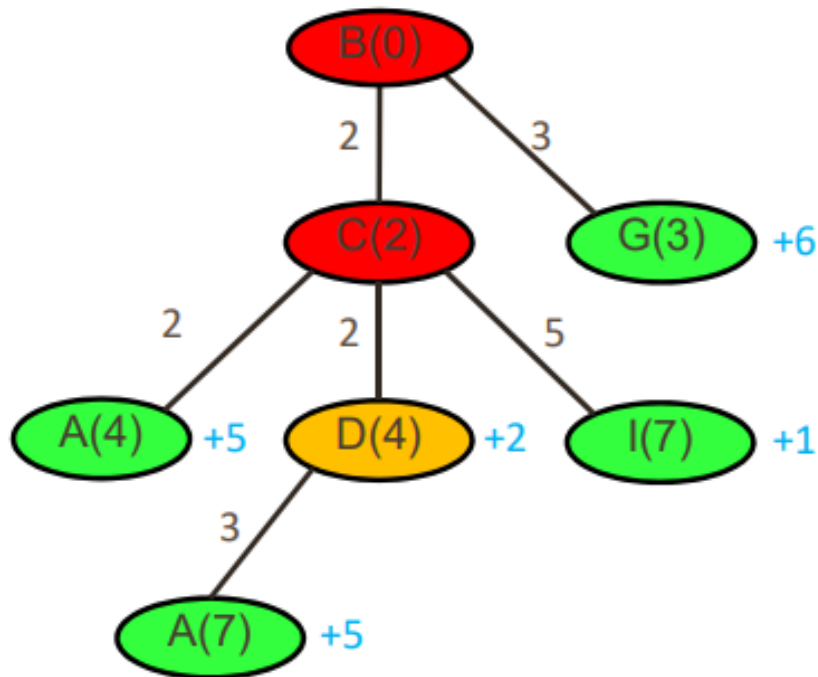
Open (Q):
{ D (4+2),
I (7+1),
G (3+5),
A (4+5) }

Closed:
{ B (0),
C (2) }





A* heuristička pretraga



Open (Q):
{ I (5+1),
G (3+5),
A (4+5) }

Closed:
{ B (0),
C (2),
D (4) }

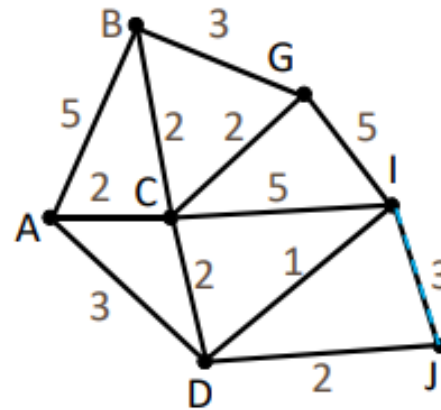
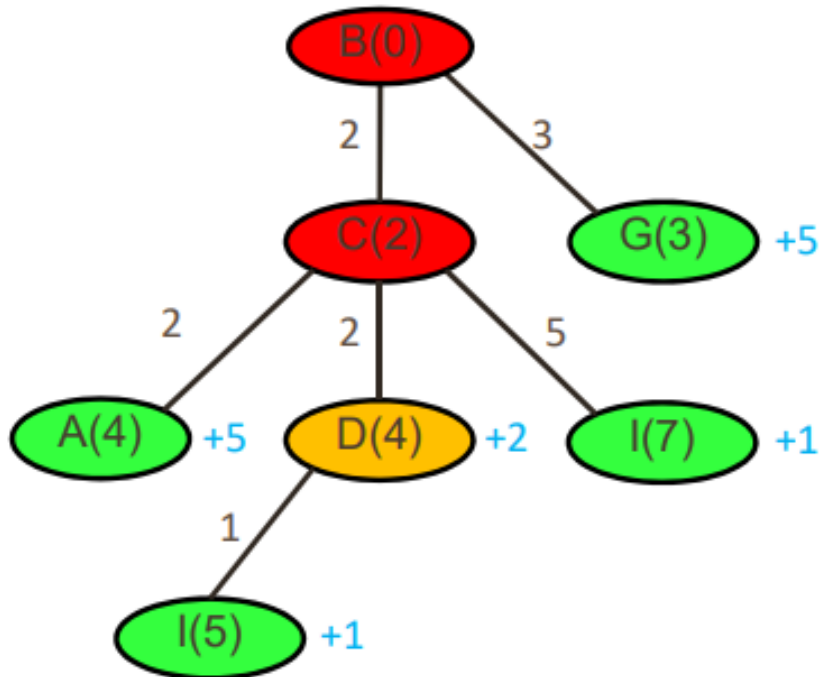




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A* heuristička pretraga



Open (Q):
{ I (5+1),
G (3+5),
A (4+5) }

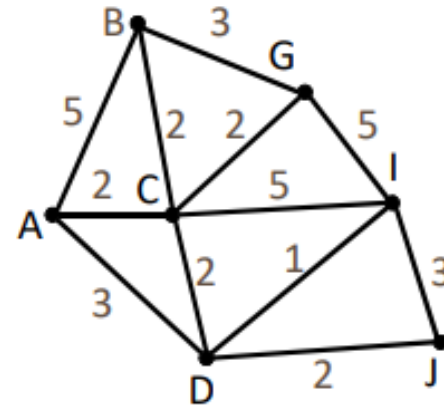
Closed:
{ B (0),
C (2),
D (4) }

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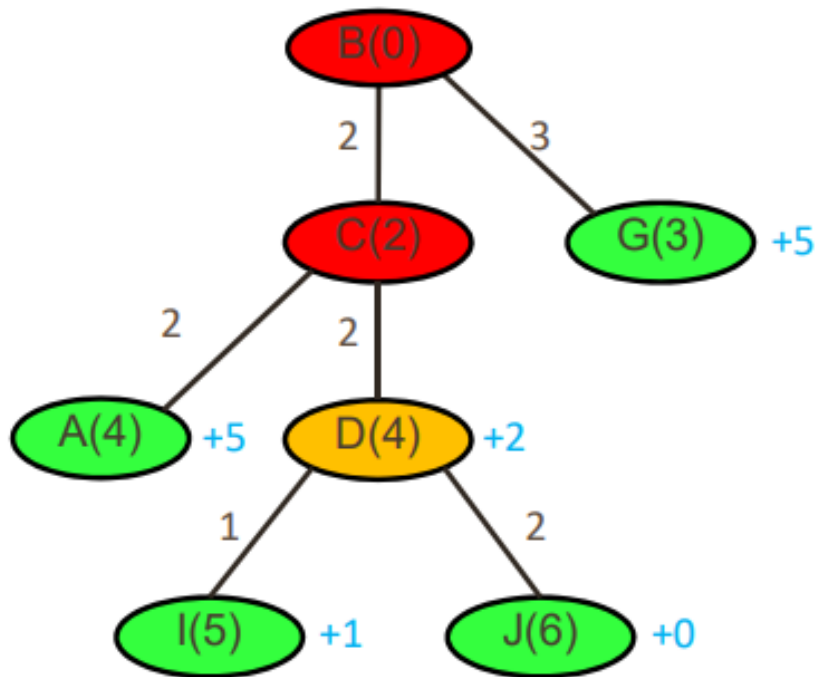


A* heuristička pretraga



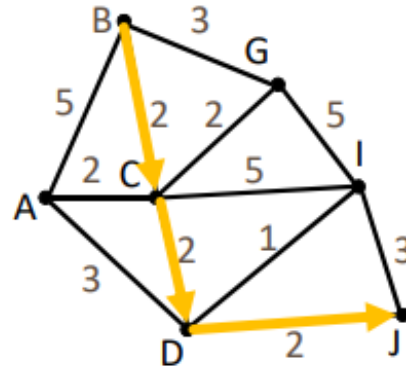
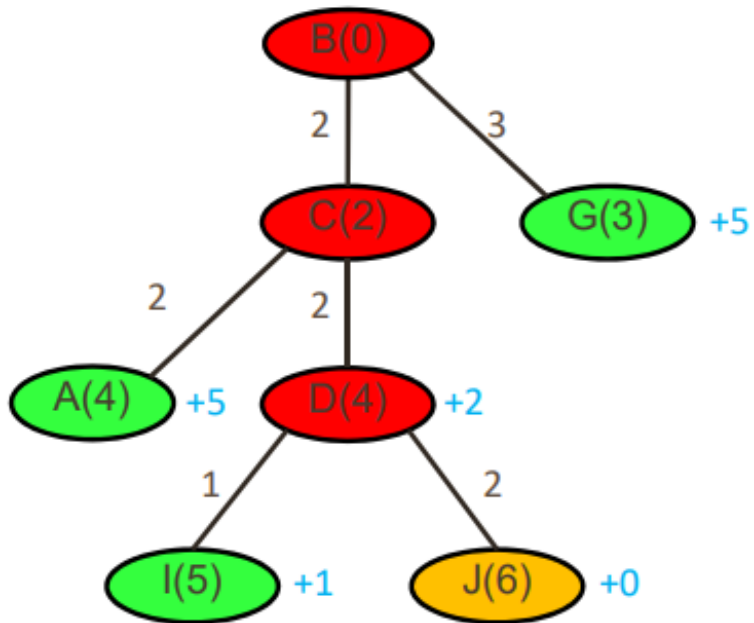
Open (Q):
{ J (6+0),
I (5+1),
G (3+5),
A (4+5) }

Closed:
{ B (0),
C (2),
D (4) }





A* heuristička pretraga



Open (Q):	Closed:
{ I (5+1),	{ B (0),
G (3+5),	C (2),
A (4+5) }	D (4),
	J (6) }

Final path solution: B→C→D→J
with path cost 6

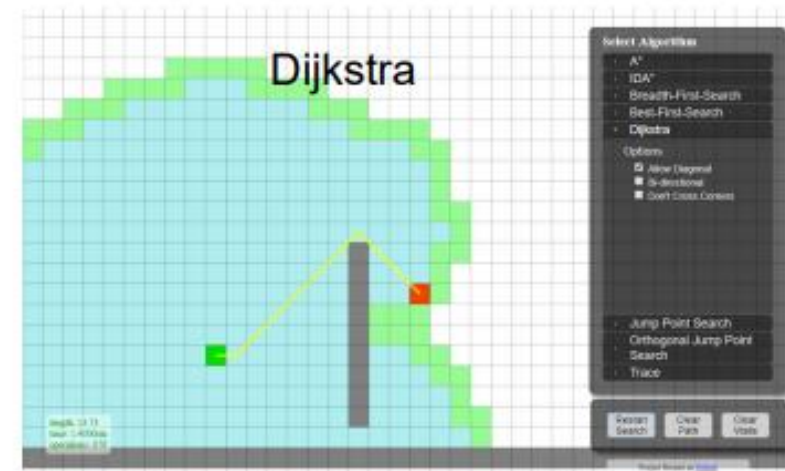
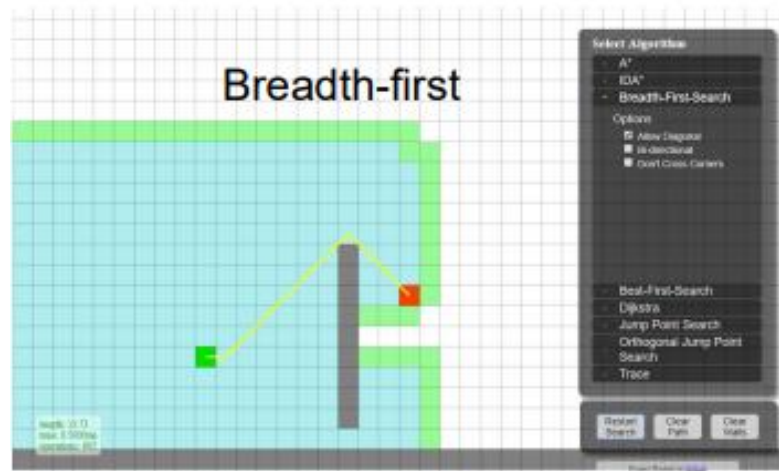




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<https://qiao.github.io/PathFinding.js/visual/>



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Thanks!

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